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Yamamoto

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(54) **PARTITIONING SORTED DATA SETS**

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G06F 7/00 (2006.01)

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CPC ***G06F 17/30312*** (2013.01); ***G06F 17/30613*** (2013.01); ***G06F 17/30864*** (2013.01)

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USPC 707/741
See application file for complete search history.

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Assistant Examiner — Raquel Perez-Arroyo
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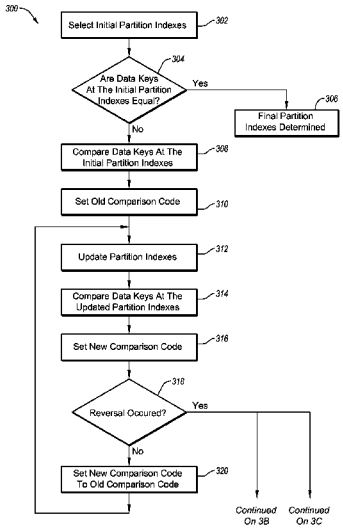
(57) **ABSTRACT**

Techniques described herein generally relate to processing data sets. The data sets can each include multiple indexed data values that can be partitioned into first and second portions. Each data value in each of the first portions of the data sets may have a greater magnitude than each data value in each of the second portions of the data sets. The first portions and second portions of the data values can be further processed by either different processors, or different processor cores as may be desirable.

22 Claims, 15 Drawing Sheets



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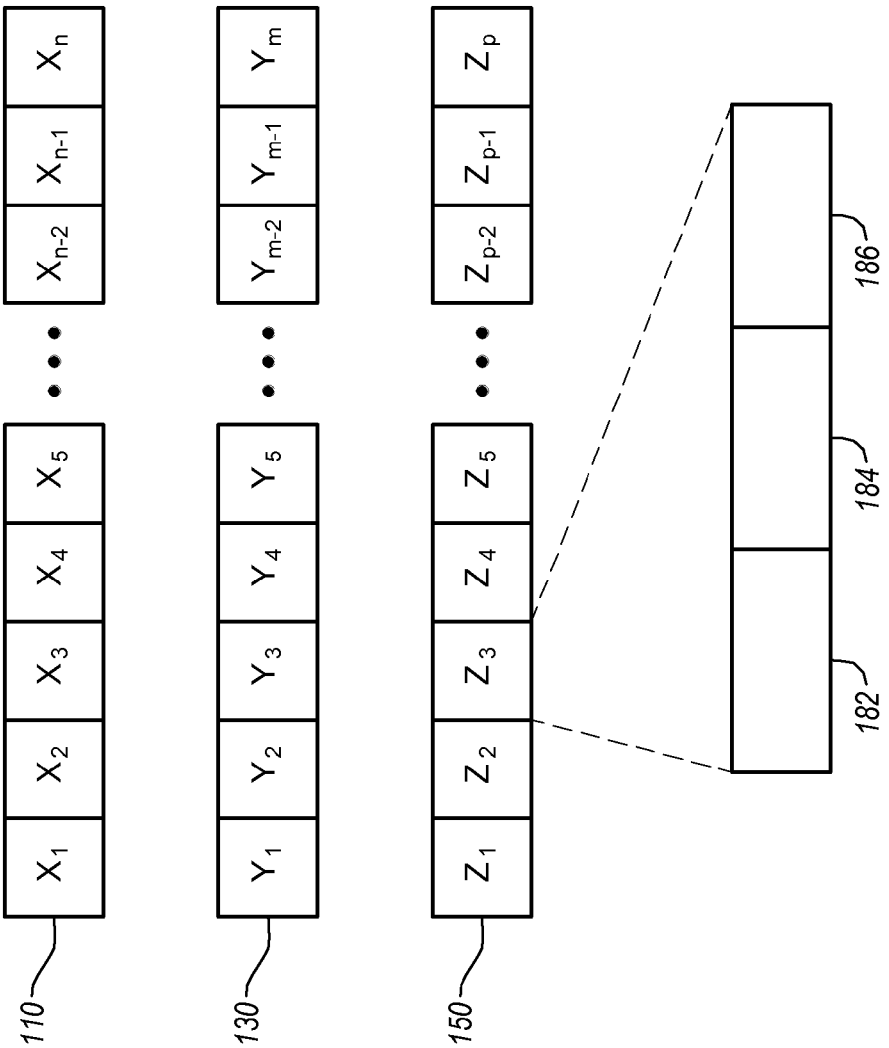
X1
X2
X3
X4
X5
Xn-2
Xn-1
Xn
110
Y1
Y2
Y3
Y4
Y5
Ym-2
Ym-1
Ym
130
Z1
Z2
Z3
Z4
Z5
Zp-2
Zp-1
Zp
150
182
184
186

Fig. 1A

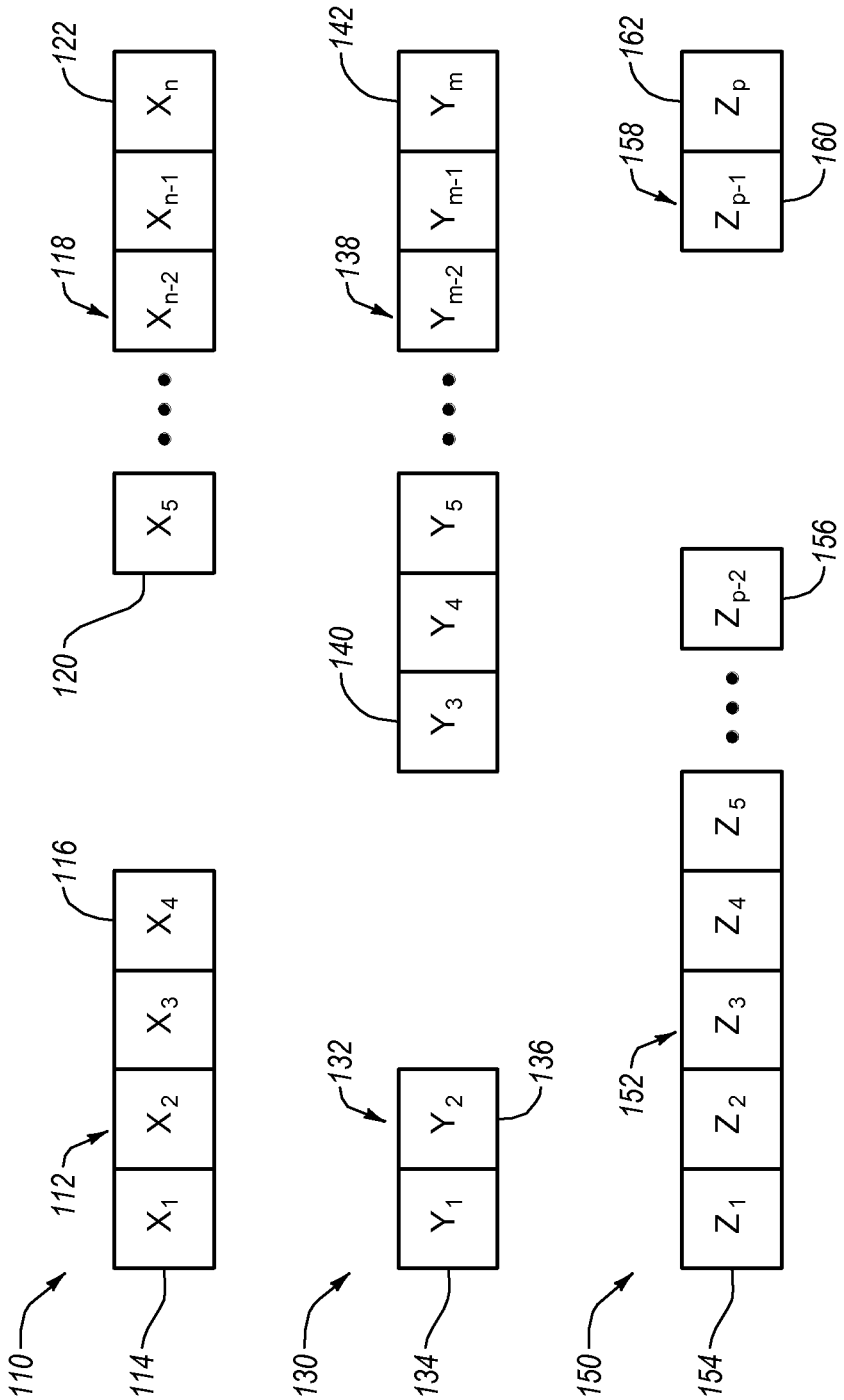
110
112
114
116
X_1
X_2
X_3
X_4
120
118
122
X_5
X_{n-2}
X_{n-1}
X_n
130
132
134
136
Y_1
Y_2
140
138
142
Y_3
Y_4
Y_5
Y_{m-2}
Y_{m-1}
Y_m
150
152
154
156
Z_1
Z_2
Z_3
Z_4
Z_5
Z_{p-2}
158
160
162
Z_{p-1}
Z_p

Fig. 1B

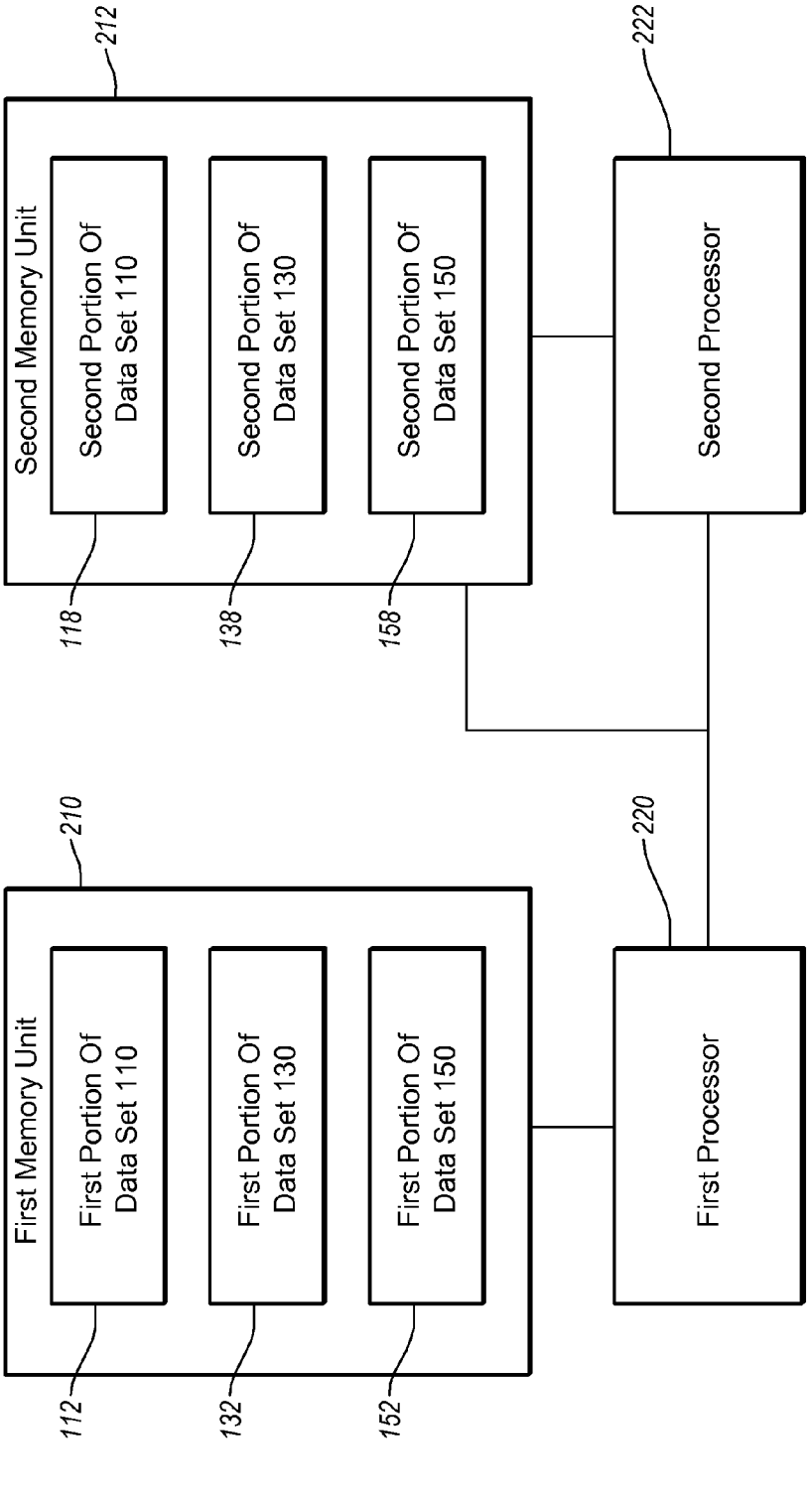
212
222
Second Memory Unit
Second Portion Of Data Set 110
Second Portion Of Data Set 130
Second Portion Of Data Set 150
Second Processor
118
138
158
210
220
First Memory Unit
First Portion Of Data Set 110
First Portion Of Data Set 130
First Portion Of Data Set 150
First Processor
112
132
152
200

Fig. 2A

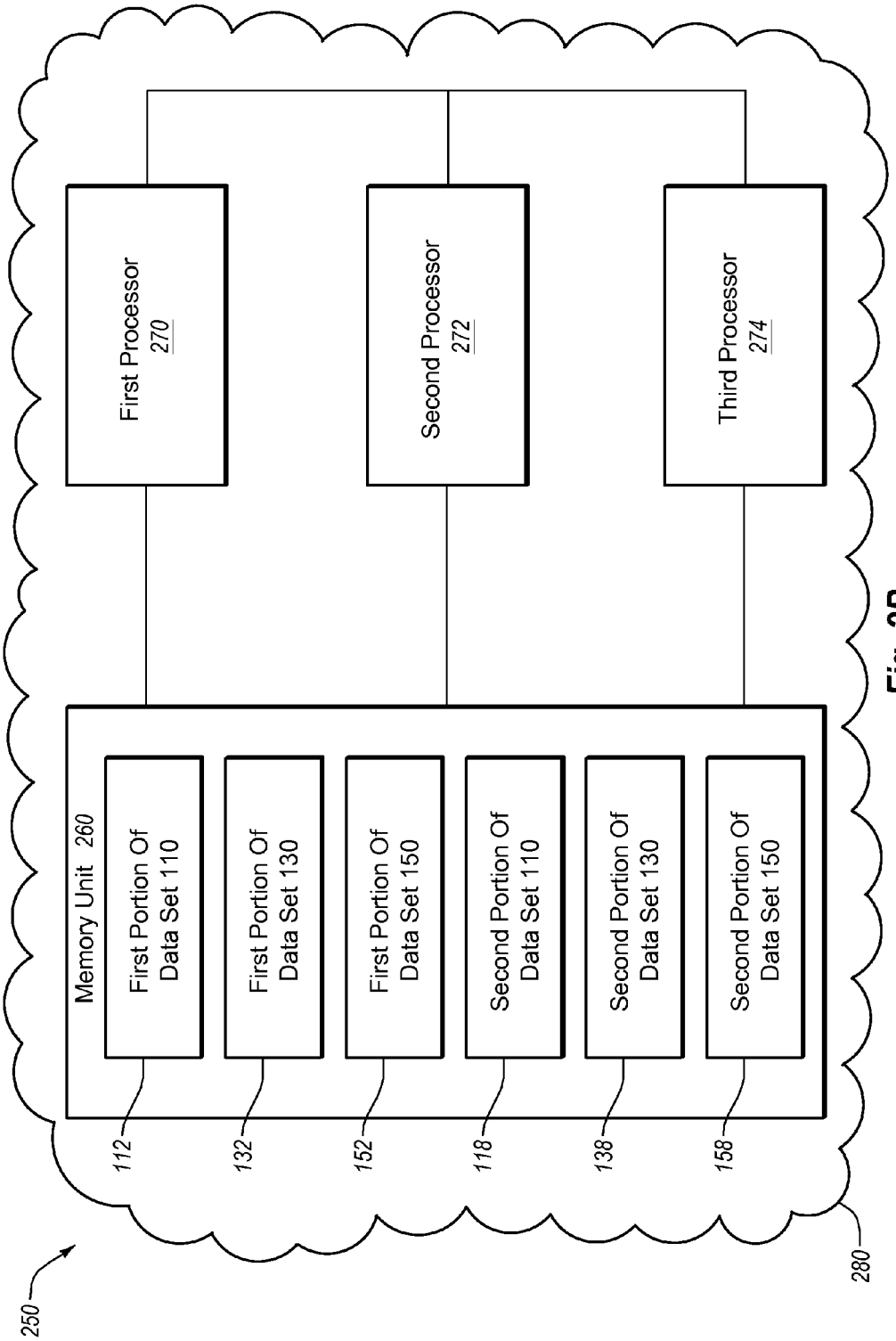
First Processor 270
Second Processor 272
Third Processor 274
Memory Unit 260
First Portion Of Data Set 110
First Portion Of Data Set 130
First Portion Of Data Set 150
Second Portion Of Data Set 110
Second Portion Of Data Set 130
Second Portion Of Data Set 150
112
132
152
118
138
158
250
280

Fig. 2B

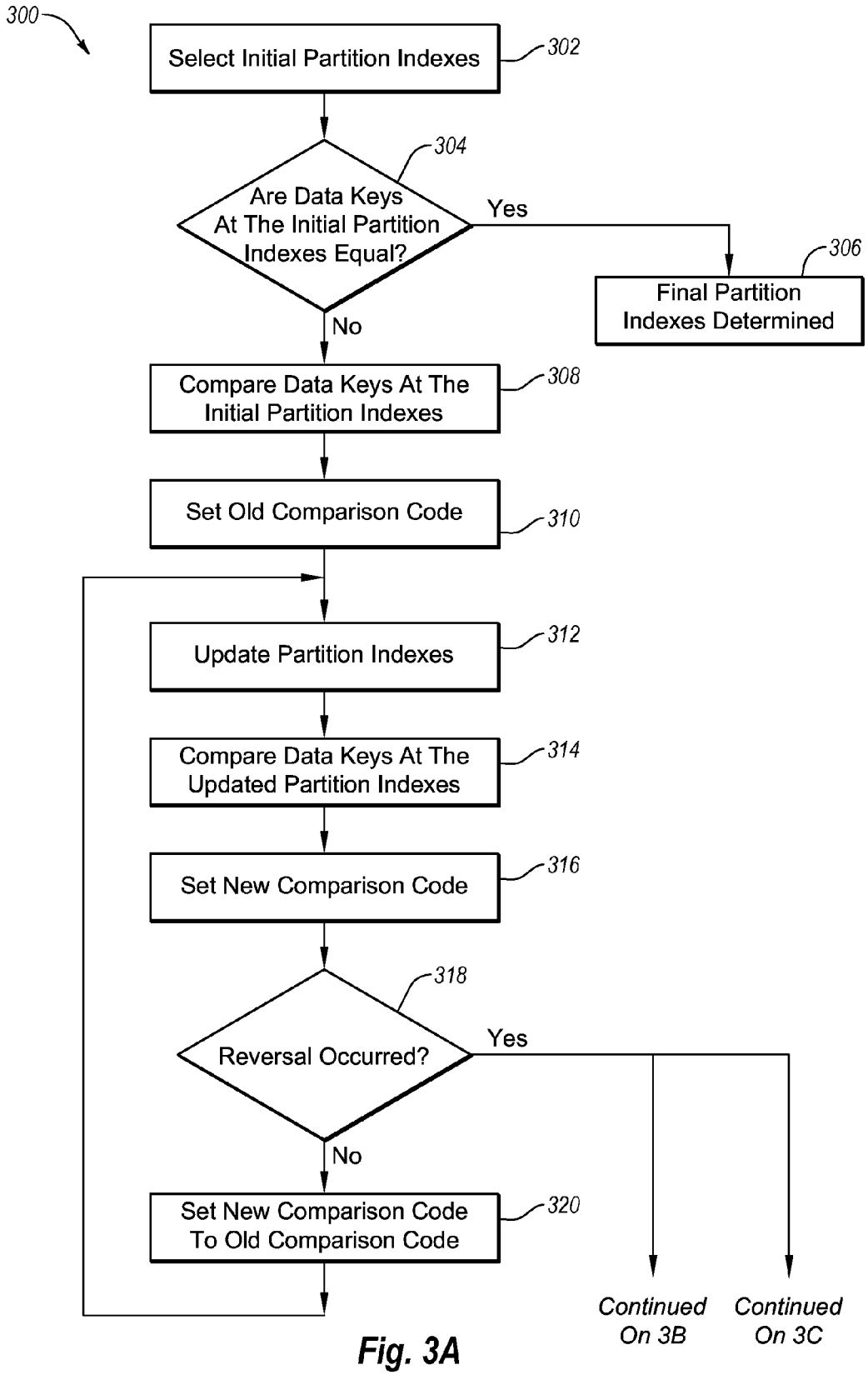
300
Select Initial Partition Indexes
302
304
Are Data Keys At The Initial Partition Indexes Equal?
Yes
306
Final Partition Indexes Determined
No
Compare Data Keys At The Initial Partition Indexes
308
Set Old Comparison Code
310
Update Partition Indexes
312
Compare Data Keys At The Updated Partition Indexes
314
Set New Comparison Code
316
318
Reversal Occurred?
Yes
No
Set New Comparison Code To Old Comparison Code
320
Continued On 3B
Continued On 3C

Fig. 3A

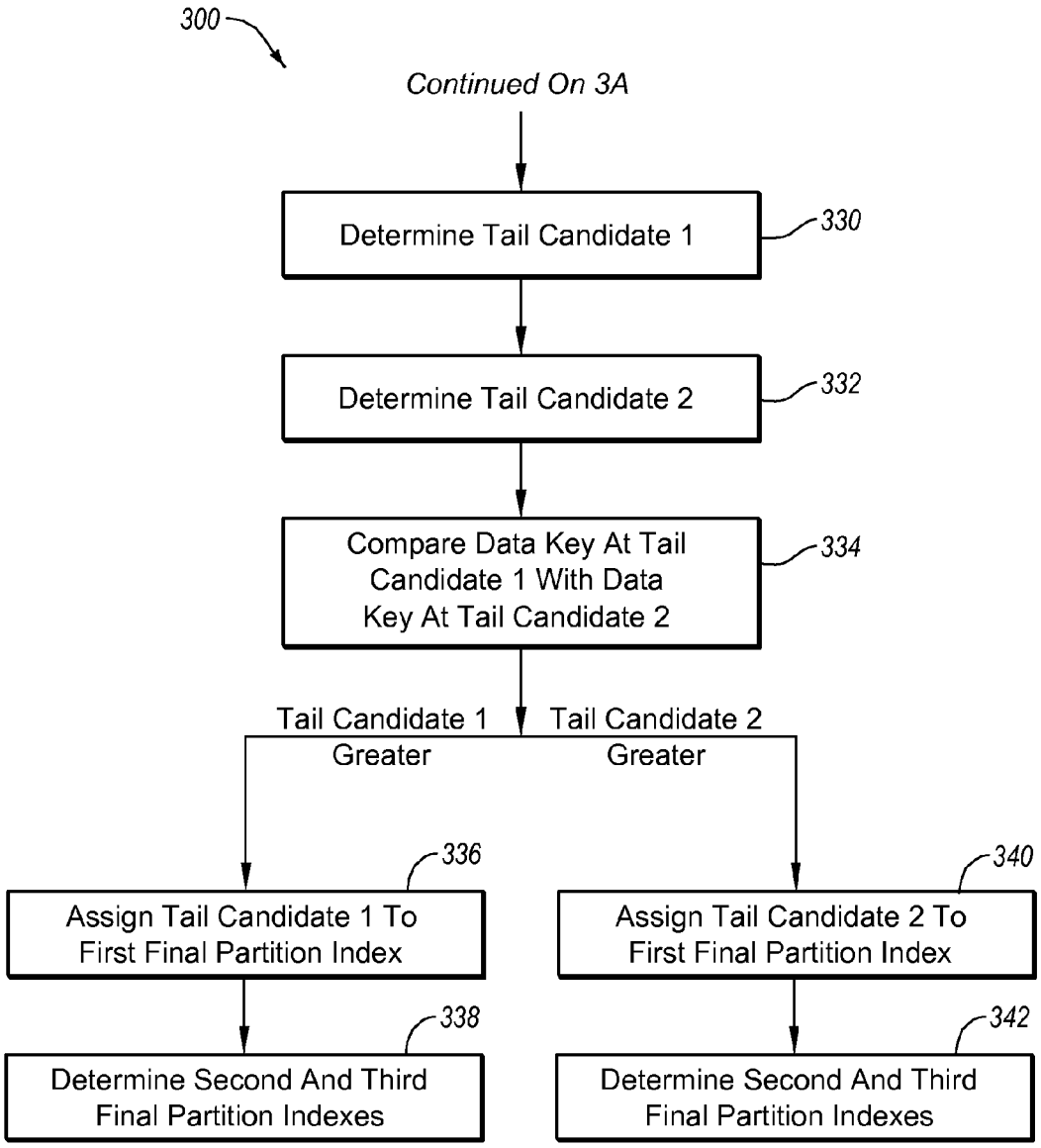
300
Continued On 3A
Determine Tail Candidate 1
330
Determine Tail Candidate 2
332
Compare Data Key At Tail Candidate 1 With Data Key At Tail Candidate 2
334
Tail Candidate 1 Greater
Tail Candidate 2 Greater
336
Assign Tail Candidate 1 To First Final Partition Index
338
Determine Second And Third Final Partition Indexes
340
Assign Tail Candidate 2 To First Final Partition Index
342
Determine Second And Third Final Partition Indexes

Fig. 3B

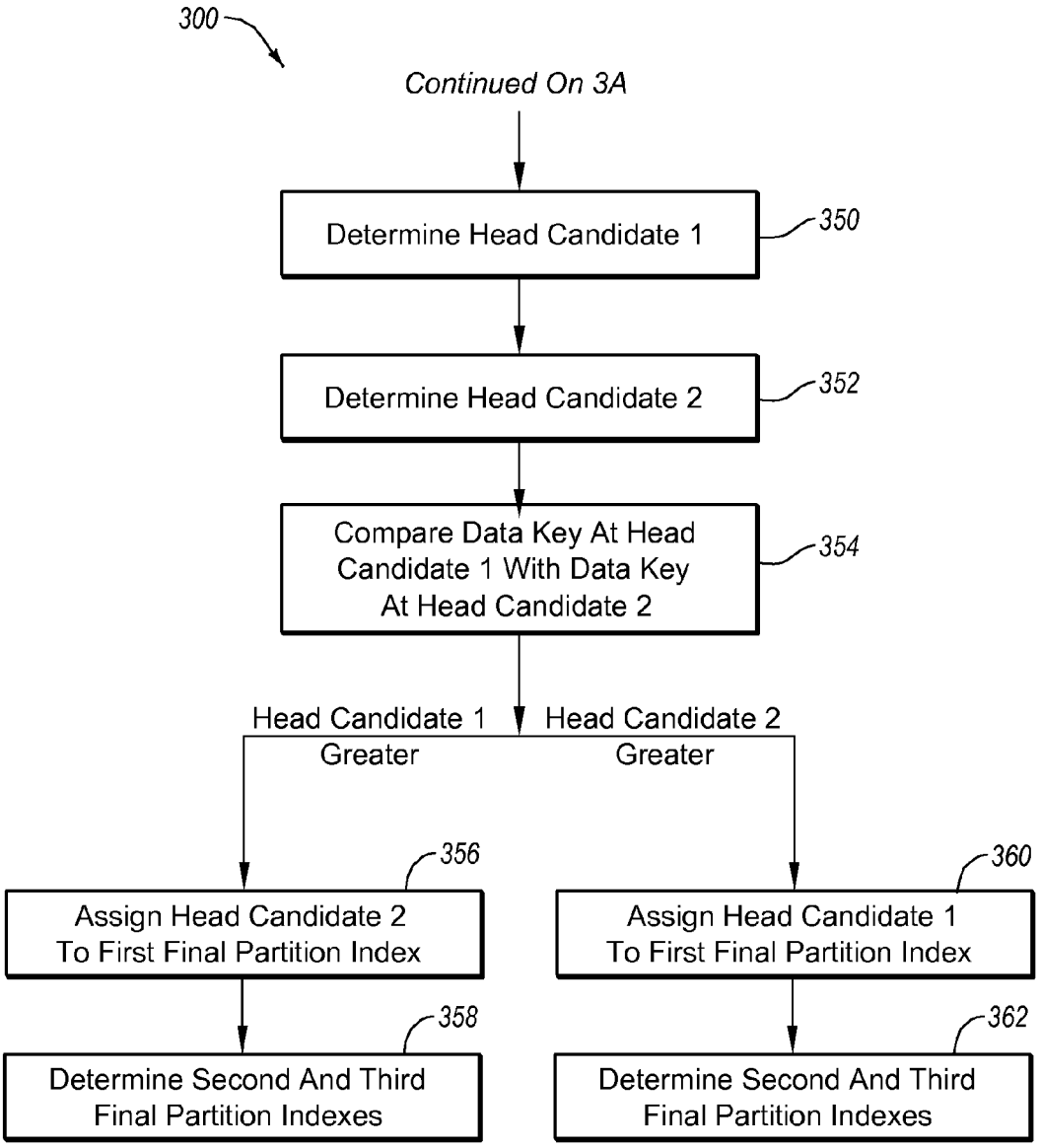
300
Continued On 3A
Determine Head Candidate 1
350
Determine Head Candidate 2
352
Compare Data Key At Head Candidate 1 With Data Key At Head Candidate 2
354
Head Candidate 1 Greater
Head Candidate 2 Greater
356
Assign Head Candidate 2 To First Final Partition Index
358
Determine Second And Third Final Partition Indexes
360
Assign Head Candidate 1 To First Final Partition Index
362
Determine Second And Third Final Partition Indexes

Fig. 3C

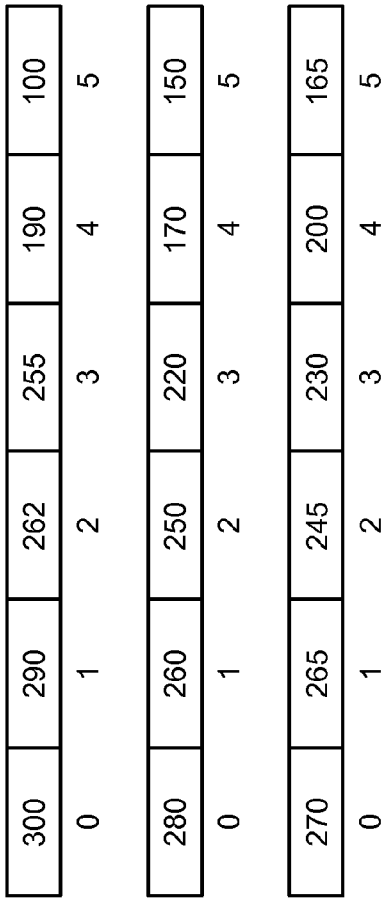
Data Set A:
300 290 262 255 190 100
Index Value i:
0 1 2 3 4 5
Data Set B:
280 260 250 220 170 150
Index Value j:
0 1 2 3 4 5
Data Set C:
270 265 245 230 200 165
Index Value k:
0 1 2 3 4 5

Fig. 4A

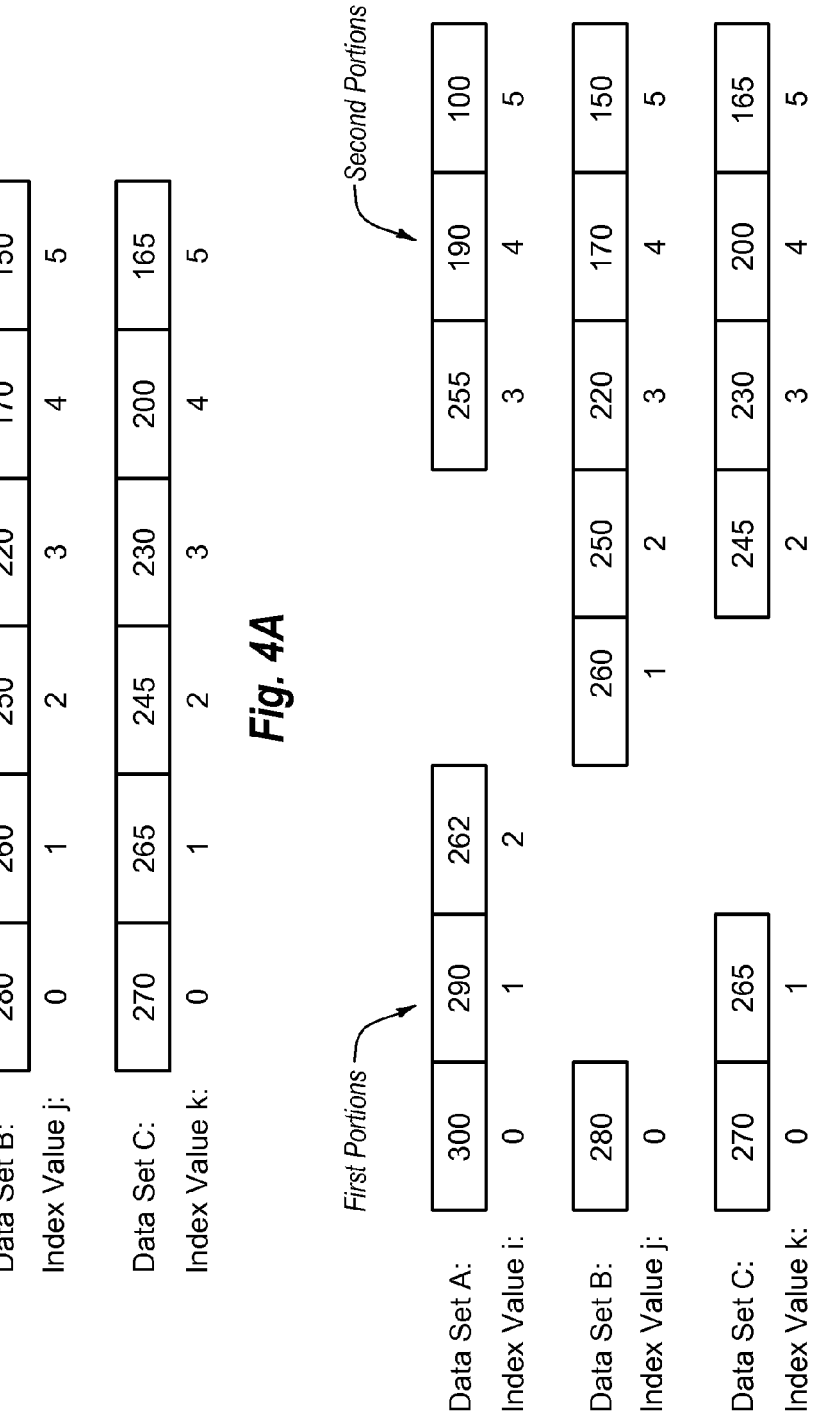
First Portions
Second Portions
Data Set A:
300 290 262
Index Value i:
0 1 2
255 190 100
3 4 5
Data Set B:
280
Index Value j:
0
260 250 220 170 150
1 2 3 4 5
Data Set C:
270 265
Index Value k:
0 1
245 230 200 165
2 3 4 5

Fig. 4B

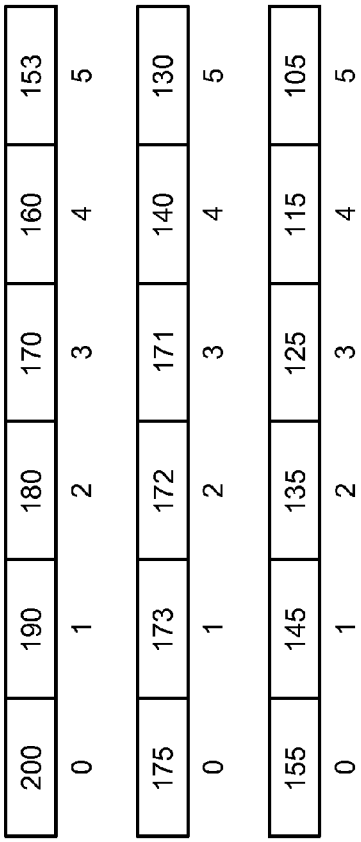
Data Set A:
200 190 180 170 160 153
Index Value i:
0 1 2 3 4 5
Data Set B:
175 173 172 171 140 130
Index Value j:
0 1 2 3 4 5
Data Set C:
155 145 135 125 115 105
Index Value k:
0 1 2 3 4 5

Fig. 5A

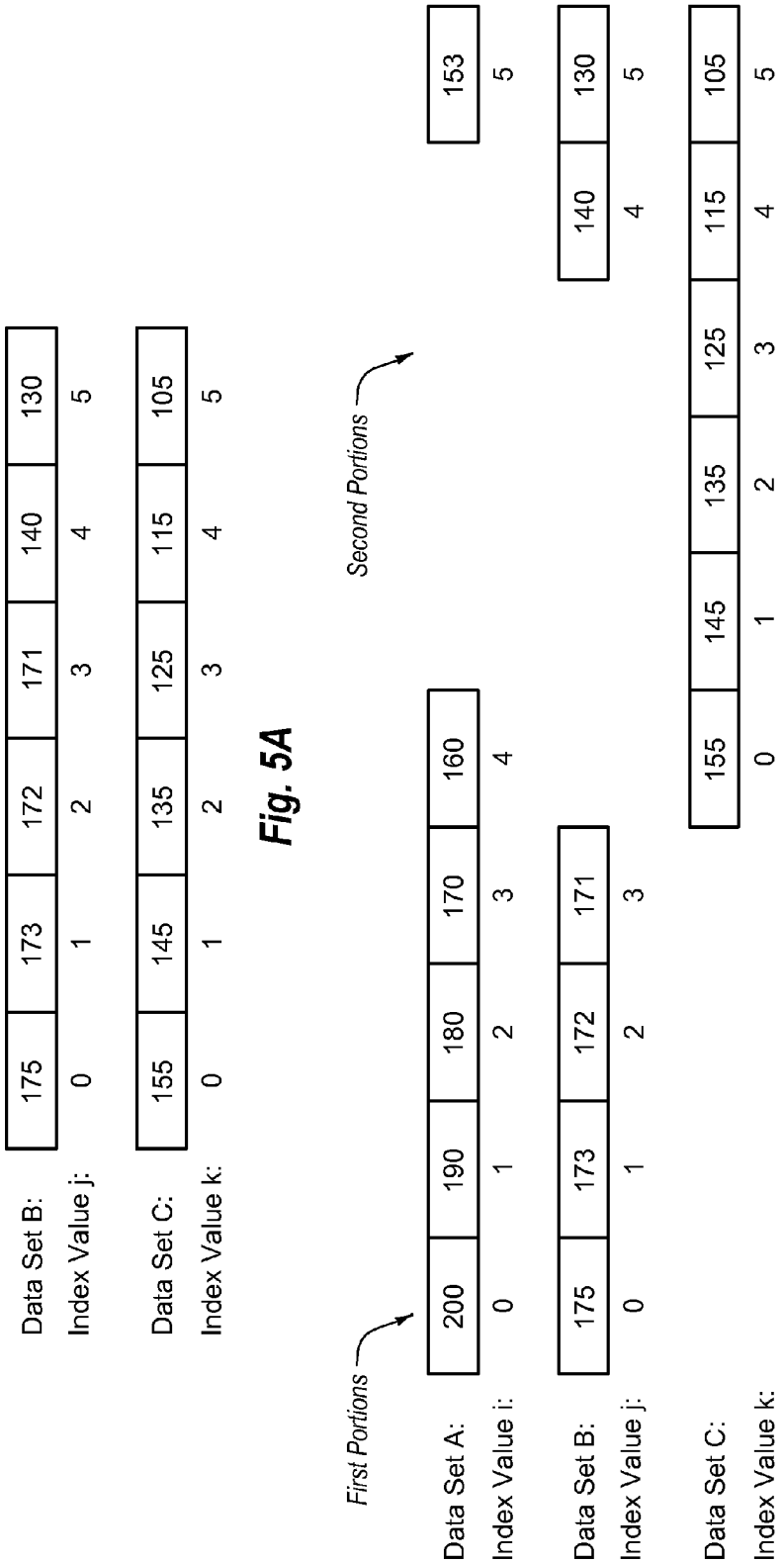
First Portions
Second Portions
Data Set A:
200 190 180 170 160
Index Value i:
0 1 2 3 4
153
5
Data Set B:
175 173 172 171
Index Value j:
0 1 2 3
140 130
4 5
Data Set C:
Index Value k:
155 145 135 125 115 105
0 1 2 3 4 5

Fig. 5B

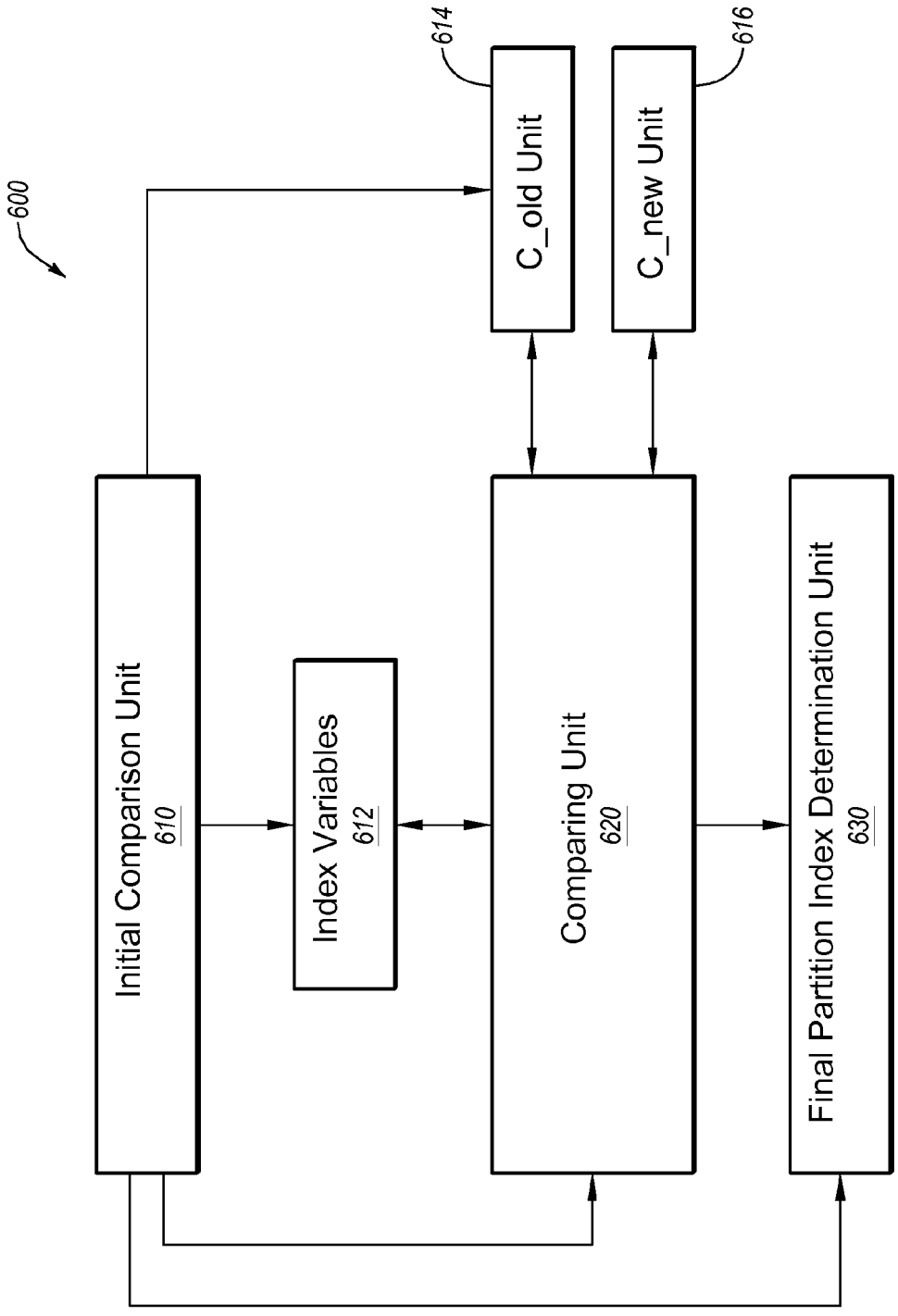
600
614
C_old Unit
616
C_new Unit
Initial Comparison Unit
610
Index Variables
612
Comparing Unit
620
Final Partition Index Determination Unit
630

Fig. 6

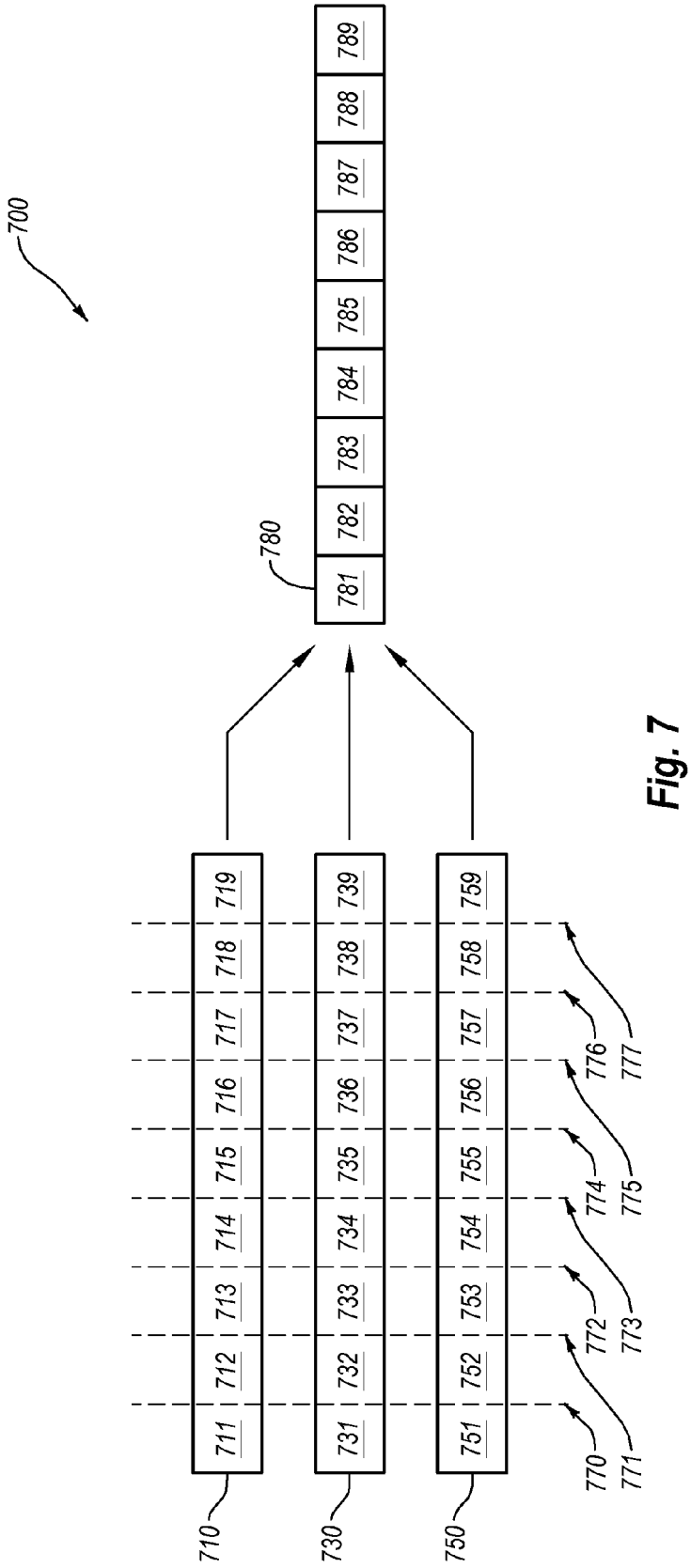
700
780
781
782
783
784
785
786
787
788
789
710
711
712
713
714
715
716
717
718
719
730
731
732
733
734
735
736
737
738
739
750
751
752
753
754
755
756
757
758
759
770
771
772
773
774
775
776
777

Fig. 7

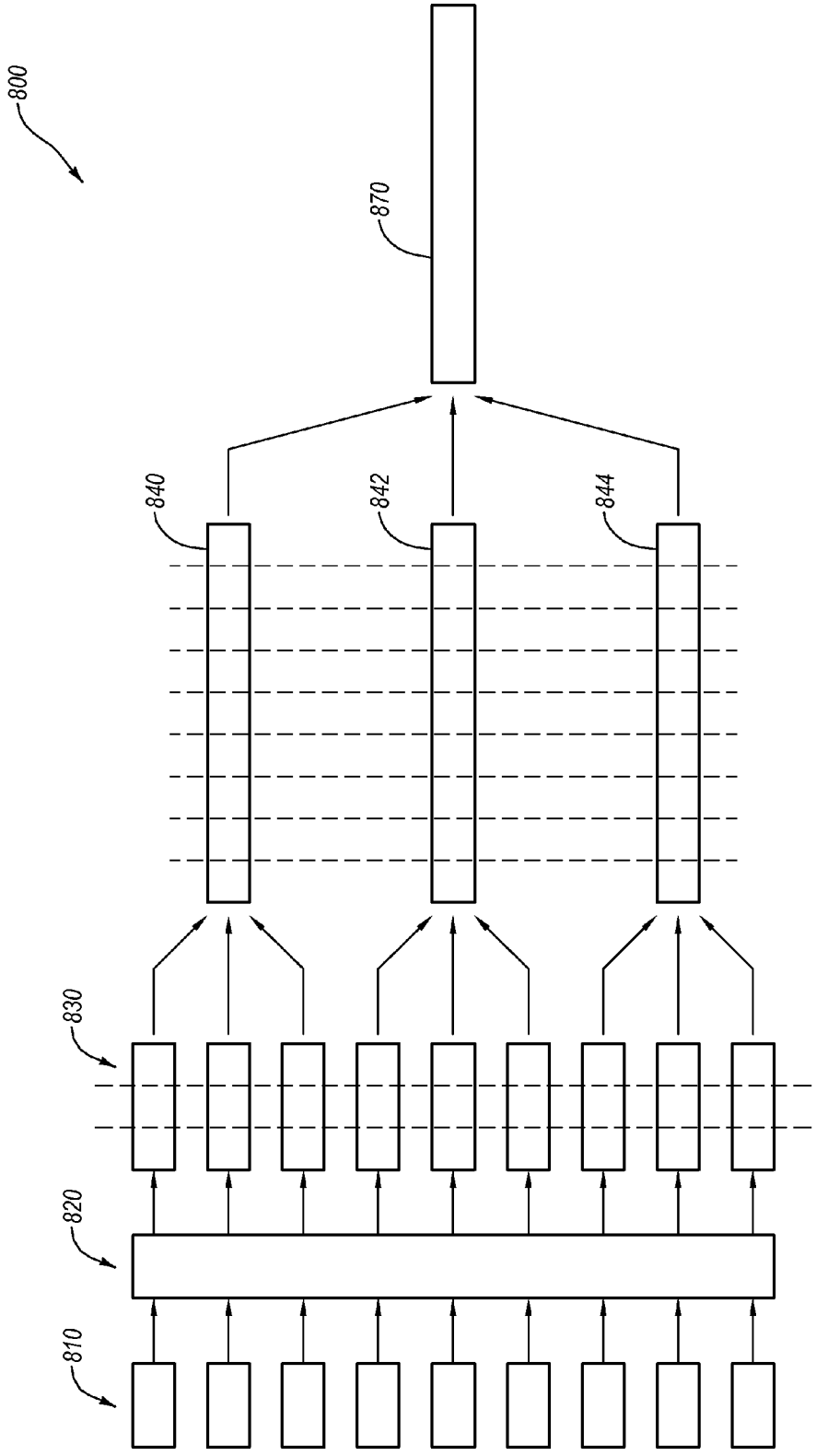

Fig. 8

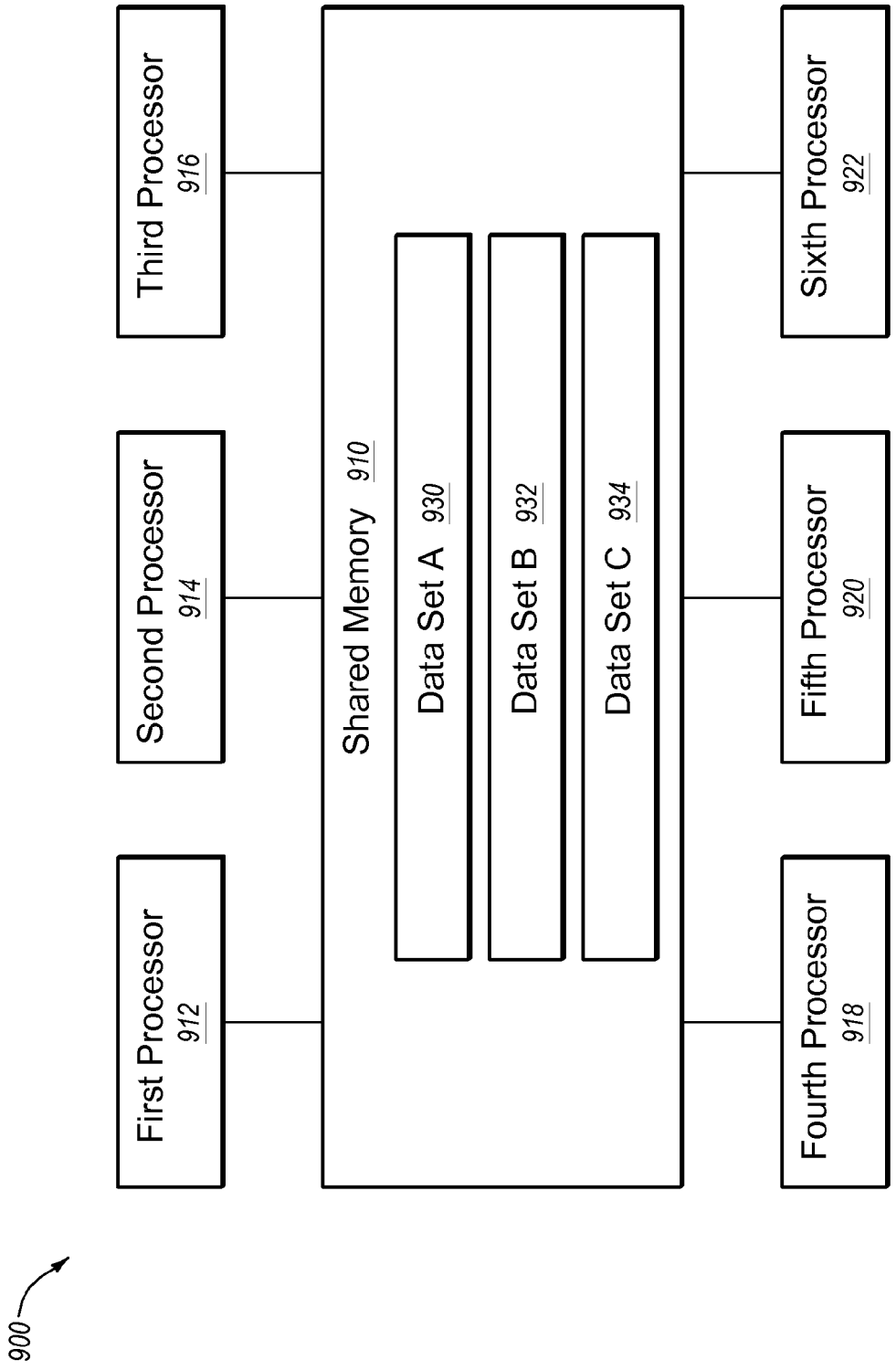
900
First Processor 912
Second Processor 914
Third Processor 916
Fourth Processor 918
Fifth Processor 920
Sixth Processor 922
Shared Memory 910
Data Set A 930
Data Set B 932
Data Set C 934

Fig. 9

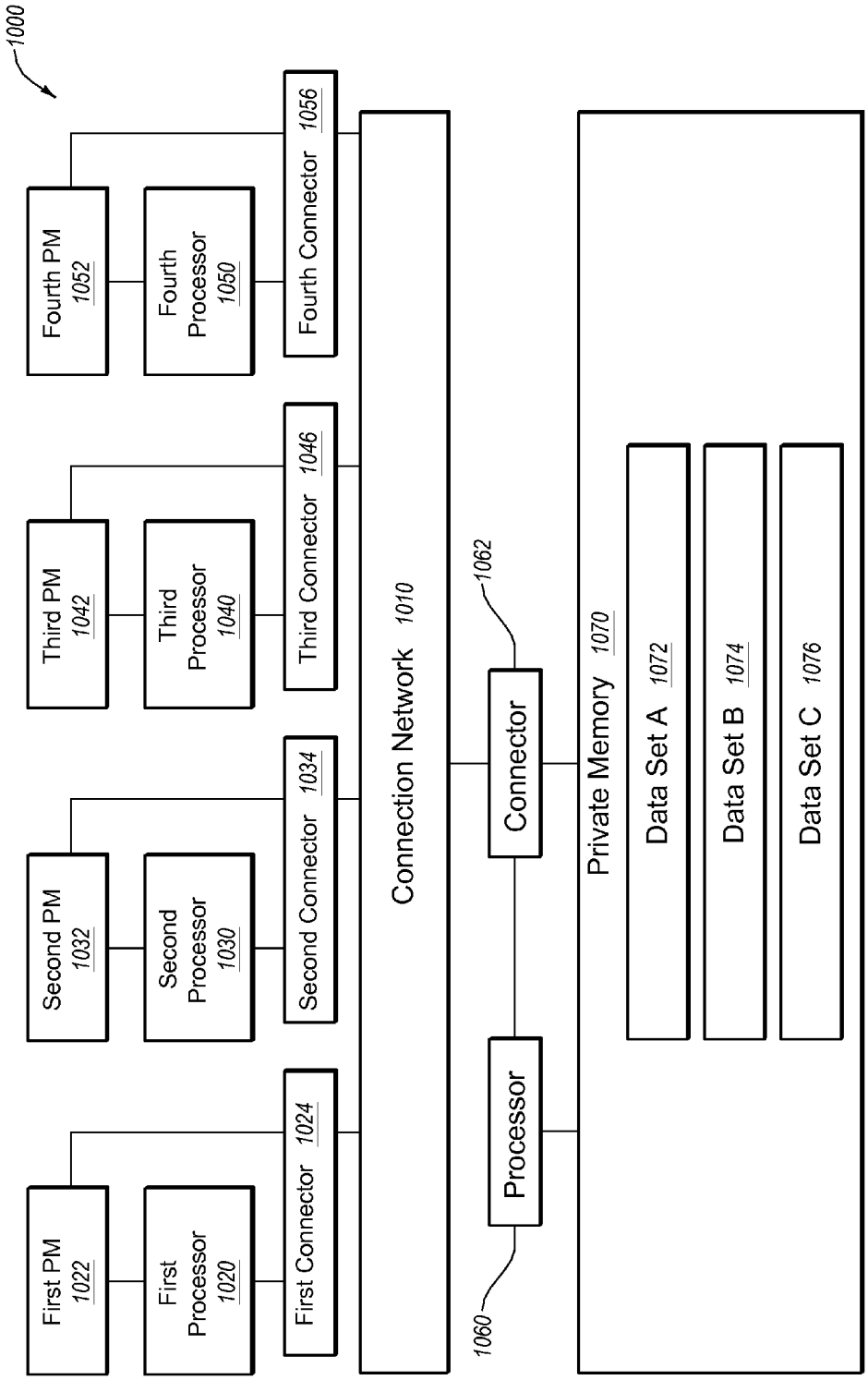
1000
Fourth PM 1052
Fourth Processor 1050
Fourth Connector 1056
Third PM 1042
Third Processor 1040
Third Connector 1046
Second PM 1032
Second Processor 1030
Second Connector 1034
First PM 1022
First Processor 1020
First Connector 1024
Connection Network 1010
Connector 1062
Processor 1060
Private Memory 1070
Data Set A 1072
Data Set B 1074
Data Set C 1076

Fig. 10

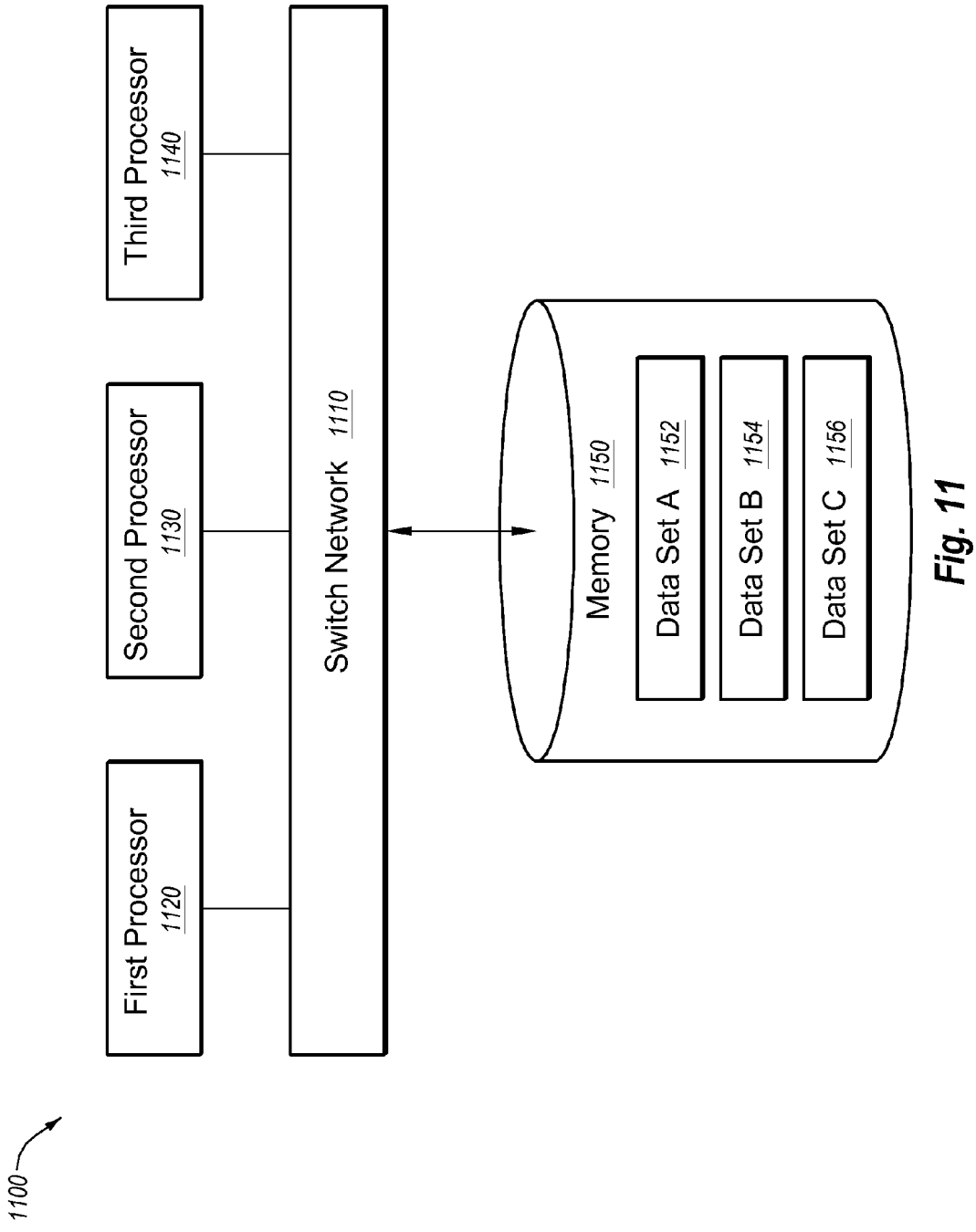
1100
First Processor 1120
Second Processor 1130
Third Processor 1140
Switch Network 1110
Memory 1150
Data Set A 1152
Data Set B 1154
Data Set C 1156

Fig. 11

PARTITIONING SORTED DATA SETS

CROSS REFERENCE TO RELATED APPLICATION

The present application is a national stage filing under 35 U.S.C. §371 of PCT Application No. PCT/US2012/051922, filed on Aug. 22, 2012.

BACKGROUND

Unless otherwise indicated herein, the materials described herein are not prior art to the claims in the present application and are not admitted to be prior art by inclusion in this section.

Information processing may include an operation of combining input (I) and memory information (S), giving processing (F) defined for an application, and acquiring output information (O). Thus, information processing may be expressed as a relation of O=F (I, S). To perform the processing, the input I and the memory information S may have a data type corresponding to the processing F. For example, if a data type is defined at a level of a Machine Instruction, a level of a compiler, and/or a level of an application, corresponding processing F may be defined for these data types.

One basic data type is an integer type. Basic operations, such as addition, subtraction, multiplication, division, and comparison may be performed at a level of machine instructions. Other operations may be performed on an integer type using a compiler and/or an application.

Another data type may be a variable-length data type, such as a character string. Various operations, such as move, copy, coupling of two or more character strings, division of a character string, searching a character string, insertion and/or deletion of the character string from a specified position may be implemented.

Another data type may be an aligned data sequence. Various operations, including but not limited to, a data sequence copy, a data sequence length calculation, a data sequence merge of two or more data sequences, a data sequence division, and a data sequence-element search may be implemented on a data sequence.

The data sequence copy operation may produce another data sequence. The data sequence length calculation operation may produce an indication of a length of a data sequence.

The data sequence division operation may divide one or more data sequences into two or more partial data sequences. The order key values within the data sequences may not be disrupted during a data sequence divide. A data sequence divide operation of a single data sequence and of 2 pair type aligned data sequences may be accomplished.

The data sequence merge operation of two or more data sequences may combine the two or more data sequences into a single data sequence. The data sequence-element search operation may determine a location of an element within a data sequence. For a 2 pair type aligned data sequence, the data sequence merge operation and the data sequence element search operations may be performed using, in part, a data sequence division operation.

In order to decrease processing time to perform some of the above operations, parallel processing of the operations may be performed. Parallel processing may be performed using multiple individual processors, a processor with multiple cores, or some combination thereof.

SUMMARY

Techniques described herein may generally relate to partitioning of data into sorted data sets.

In some examples, a method for a computing device is described that may include locating a first partition index for each of three sorted data sets where each sorted data set may include multiple indexed data values. Each of the first partition indexes may identify an index location that may be utilized to partition the corresponding data set into first and second portions of data values. Each data value in each of the first portions of the data sets may have a greater magnitude than each data value in each of the second portions of the data sets. Locating the partition index for each of the data sets may include selecting an initial partition index for each data set and comparing the data values at the initial partition index for each data set to identify a highest data value and a lowest data value. Locating the partition index may also include adjusting the initial partition index for the data sets with the highest and lowest data values and comparing data values at the adjusted partition indexes of the two data sets and the data value at the initial partition index of the data set without an adjusted partition index.

In some examples, a method is described that may include partitioning, by a computing device, three sorted data sets that each include multiple indexed data values into first and second portions. Each data value in each of the first portions of the sorted data sets may have a greater magnitude than each data value in each of the second portions of the sorted data sets. The method may also include processing the first portions of the sorted data sets. The method may also include processing the second portions of the sorted data sets.

In some examples, a system is described that may include a first processor configured to process a first portion of each of three sorted data sets and a second processor configured to process a second portion of each of the three sorted data sets. The three sorted data sets may each include multiple indexed data values. The three sorted data sets may be partitioned into the first and second portions so that each data value in each of the first portions of the sorted data sets may have a greater magnitude than each data value in any of the second portions of the sorted data sets. The first and second processor may be either two separate processors, or two processor cores within the same processor.

In some examples, a computer-readable storage medium is described whose contents, when executed by a processor, may cause the processor to locate a first partition index for each of three sorted data sets. Each data set may include multiple indexed data values. Each first partition index may identify an index location that may be utilized to partition each corresponding sorted data set into a first portion of data values and a second portion of data values. Each data value in each of the first portions of the sorted data sets may have a greater magnitude than each data value in each of the second portions of the sorted data sets. Execution of the contents may also cause the processor to select an initial partition index for each sorted data set. Execution of the contents may also cause the processor to compare the data values at the initial partition index for each sorted data set to identify a highest data value and a lowest data value. Execution of the contents may also cause the processor to adjust the initial partition index for the sorted data set with the highest data value and for the sorted data set with the lowest data value. Execution of the contents may also cause the processor to compare data values at the adjusted partition indexes of the two data sets having the adjusted partition indexes and the data value at the initial partition index of the sorted data set without an adjusted partition index.

The foregoing summary is illustrative only and is not intended to be in any way limiting. In addition to the illustrative aspects, embodiments, and features described above, fur-

ther aspects, embodiments, and features will become apparent by reference to the drawings and the following detailed description.

BRIEF DESCRIPTION OF THE FIGURES

In the drawings:

FIG. **1**A illustrates three example data sets;

FIG. **1**B illustrates an example partitioning of the data sets illustrated in FIG. **1**A;

FIG. **2**A illustrates an example system configured to process the partitioned portions of the data sets of FIG. **1**B;

FIG. **2**B illustrates another example system configured to process the partitioned portions of the data sets of FIG. **1**B;

FIGS. **3**A, **3**B, and **3**C show an example flow diagram of a method for determining a partition index for each of three sorted data sets;

FIG. **4**A illustrates three example data sets;

FIG. **4**B illustrates an example partitioning of the data sets illustrated in FIG. **4**A

FIG. **5**A illustrates three example data sets;

FIG. **5**B illustrates an example partitioning of the data sets illustrated in FIG. **5**A;

FIG. **6** illustrates an example system for implementing the method of FIGS. **3**A-**3**C;

FIG. **7** illustrates an example merge of three example data sets;

FIG. **8** illustrates an example merge of nine example data sets;

FIG. **9** illustrates an example system for partitioning data sets;

FIG. **10** illustrates another example system for partitioning data sets; and

FIG. **11** illustrates another example system for partitioning data sets, all arranged in accordance with at least some embodiments described herein.

DETAILED DESCRIPTION

In the following detailed description, reference is made to the accompanying drawings, which form a part hereof. In the drawings, similar symbols typically identify similar components, unless context dictates otherwise. The illustrative embodiments described in the detailed description, drawings, and claims are not meant to be limiting. Other embodiments may be utilized, and other changes may be made, without departing from the spirit or scope of the subject matter presented herein. It will be readily understood that the aspects of the present disclosure, as generally described herein, and illustrated in the Figures, can be arranged, substituted, combined, separated, and designed in a wide variety of different configurations, all of which are explicitly contemplated herein.

Some embodiments described herein may generally relate to information processing of aligned data sets, such as a merge operation, a division operation, a data element searching operation, among other operations that may be performed on aligned data sets. In particular, some embodiments describe a method of partitioning 3-pair type aligned data sets into an arbitrary number of partial data sets, without breaking the order of alignment of data keys used to sort the data sets. Partitioning data sets into partial data sets may allow for greater ease in information processing of the data sets. For example, partitioning 3-pair type aligned data sets may allow for a merger of the 3-pair type aligned data sets. Additionally, the method of partitioning the 3-pair type aligned data sets as described herein may be performed using parallel processing thereby decreasing an amount of time from start to finish to perform the division.

A general method to perform the division of 3-pair type aligned data sets may be as follows. Three sorted data sets may each be partitioned into first and second portions, where each value in each of the first portions of each data set may have a greater magnitude than each value in each of the second portions of each data set. Partitioning the three data sets into the first and second portions may allow parallel processing of the first and second portions of the three data sets. For example, the first portions of each of the three data sets may be merged and sorted by a first processor and the second portions of each of the three data sets may be merged and sorted by a second processor. The merged and sorted first and second portions may also be merged together. The merged and sorted first and second portions may be equivalent to the result of merging and then sorting the three data sets. Furthermore, in some embodiments, the three data sets may be independently partitioned into more than two portions. Each of the portions may be processed individually and in parallel and the results combined.

FIG. **1**A illustrates three example data sets, arranged in accordance with at least some embodiments described herein. The three example data sets are described generally herein as data sets that may be partitioned based on the disclosure contained herein to generate partial data sets of an arbitrary number without disrupting the order of alignment of data keys within data elements that form the three example data sets.

The illustrated embodiment may include an x data set **110**, a y data set **130**, and a z data set **150**. The x data set **110** may include n number of indexed data elements $X_1, X_2 \ldots X_n$. The y data set **130** may include m number of indexed data elements $Y_1, Y_2 \ldots Y_m$. The z data set **150** may include p number of indexed data elements $Z_1, Z_2 \ldots Z_p$. Each data element in each of the data sets **110**, **130**, **150** may include one or more data values. For example, FIG. **1**A illustrates that the data element Z_3 may include first, second, and third data values **182**, **184**, **186**. Each of the data values **182**, **184**, **186** may be a different type of data. For example, in some embodiments, data value **182** may be an identification number, data value **184** may be a name, and data value **186** may be an age.

Any type of data or combination of data may be included in the data elements of the data sets **110**, **130**, **150**. Each data element of each data set **110**, **130**, **150** may contain at least one type of data value that is sortable. That is, each data element of each data set **110**, **130**, **150** may contain at least one type of data value that may be organized in a logical order, such as, ascending or descending order. For example, the sortable data values may be numbers or words that may be sorted in a logical order such as ascending or descending numerical order, ascending or descending alphabetical order, or the like.

The data elements of each of the data sets **110**, **130**, **150** may be sorted by at least one of the sortable data values within the data element. The data value that is used to sort the data elements within a data set may be referred to herein as a data key. The data elements of each of the data sets **110**, **130**, **150** may each contain the same type of data key by which the data sets **110**, **130**, **150** may be sorted. In some embodiments, the data sets **110**, **130**, **150** may contain multiple data values that may be data keys. In these and other embodiments, the data sets **110**, **130**, **150** may be aligned using one of the multiple data keys in either ascending or descending order.

As an example, the data sets **110**, **130**, **150** may be any sorted data sets that contain information about customers and individuals used by credit card companies, banking compa-

nies, e-commerce management companies, utility companies, such as cell phone companies, power companies, gas companies, and water supply companies, as well as other companies. In these and other examples, a data element within the data sets **110**, **130**, **150** may be a group or collection of data values that represent information about an individual that is a customer, potential customer, or someone of interest to the company that maintains the sorted data sets **110**, **130**, **150**. As one example, the data values within the data element may include an identification number assigned to an individual, the name of the individual, and the credit score of the individual.

As another example, the data sets **110**, **130**, **150** may be any sorted data sets that contain information used in databases for managing inventory of merchandise and other products. In these and other examples, a data element within the data sets **110**, **130**, **150** may be a group or collection of data values that represent information about a particular type of product of a company that maintains the sorted data sets **110**, **130**, **150**. As one example, the data values within the data element may include a product identification number, such as a UPC bar code number, the total number of products in inventory, the locations where the product is located, and other information about the product. The foregoing examples of data sets are given only as examples and are not limiting in any way.

In some embodiments, the number of data elements n, m, and p within each data set **110**, **130**, **150**, respectively may be the same. In other embodiments, the number of data elements n, m, and p within each data set **110**, **130**, **150**, respectively, may be different. In some embodiments, the data elements in the data sets **110**, **130**, **150**, may include one, two, three, four, or more data values. In some embodiments, the data elements in one of the data sets **110**, **130**, **150** may contain more data values than the data elements in another of the data sets **110**, **130**, **150**. Alternately or additionally, the data elements within each data set **110**, **130**, **150** may contain the same number of data values. Furthermore, the data elements within the data sets **110**, **130**, **150** may be indexed. For example, the first data element of the data set **110** may have an index value of 0 and the last data element of the data set **110** may have an index value of n−1, with the remaining data elements having corresponding index values between 0 and n−1.

The data sets **110**, **130**, **150** may be used in or subjected to various processing operations. For example, the data sets **110**, **130**, **150** may be merged or have some other processing operation performed thereon that generates one or more new data sets. Alternately or additionally, processing operations such as element retrieval and other non-altering processing operations may be performed on the data sets **110**, **130**, **150**.

The present disclosure provides for partitioning the data sets **110**, **130**, **150** and/or any 3-pair type aligned data sets into the 3-pair type aligned partial data sets of an arbitrary number without disrupting the order of alignment of data keys of the data sets. Partitioning the data sets **110**, **130**, **150** may assist in the merging, element retrieval or other operation(s) performed on the data sets **110**, **130**, **150**.

The data sets **110**, **130**, **150** may be partitioned into two or more portions. In particular, the data sets **110**, **130**, **150** may be partitioned into two or more portions where the data keys used to sort the data elements in a first portion of the data sets **110**, **130**, **150** are greater in magnitude than the data keys used to sort data elements in a second portion of the data sets **110**, **130**, **150**. For example, the range of sorted data keys in a first portion may be from 100 to 50 and the range of sorted data keys in a second portion may be from 48 to 5.

Partitioning of the data sets **110**, **130**, **150** may not be limited to sorted data keys with numerical values (e.g., floating point values or integer values). For example, the data sets **110**, **130**, **150** may be sorted by data keys that contain words or characters. In the logical order assigned to the sorted data keys, a first portion of the data sets **110**, **130**, **150** may contain data keys of greater magnitude than data keys in the second portion of the data sets **110**, **130**, **150**. For example, the first portion may contain sorted data keys, such as words, that start with the letters A to M and the second portion may contain data keys that start with the letters N to Z. Thus, the partitioning of the data sets **110**, **130**, **150** may maintain the logical order of the sorted data keys so that no data key in a first portion would logically fall between data keys in a second portion.

FIG. **1**B illustrates an example partitioning of the data sets **110**, **130**, **150** illustrated in FIG. **1**A, arranged in accordance with at least some embodiments described herein. The x data set **110** may be partitioned into first and second portions **112**, **118**. The first portion **112** may have a head data element **114** and a tail data element **116**. In general, a head data element may be a first data element in a portion of a data set and a tail data element may be a last data element in a portion of a data set. Thus, the second portion **118** may also have a head data element **120** and a tail data element **122**. The tail data element **116** of the first portion **112** may be the data element that was adjacent to the head data element **120** of the second portion **118** before the x data set **110** was partitioned.

The y data set **130** may be partitioned into first and second portions **132**, **138**. The first portion **132** may have a head data element **134** and a tail data element **136**. The second portion **138** may have a head data element **140** and a tail data element **142**. The tail data element **136** of the first portion **132** may be the data element that was adjacent to the head data element **140** of the second portion **138** before the y data set **130** was partitioned.

The z data set **150** may be partitioned into first and second portions **152**, **158**. The first portion **152** may have a head data element **154** and a tail data element **156**. The second portion **158** may have a head data element **160** and a tail data element **162**. The tail data element **156** of the first portion **152** may be the data element that was adjacent to the head data element **160** of the second portion **158** before the z data set **150** was partitioned.

The division of the data sets **110**, **130**, **150** may satisfy two conditions. The first condition may be the first portions **112**, **132**, **152** and the second portions **118**, **138**, **158** having tail and head data elements respectively, with the minimum value of the data keys in the tail data elements in the first portions **112**, **132**, **152** being larger or equal to the maximum value of the data keys in the head data elements in the second portions **118**, **138**, **158** as illustrated in the following equation:

min[tail data element 116,tail data element 136,tail data element 156]>=max[head data element 120, head data element 140,head data element 160].

The second condition may be that the number of data elements in the first portions **112**, **132**, **152** is equal to 3X where X is a value obtained by dividing the total number of data elements in the first portions **112**, **132**, **152** by 3.

The first portions **112**, **132**, **152** having data keys in the tail data elements that are larger in magnitude than the data keys in the head data elements in the second portions **118**, **138**, **158** may be referred to as maintaining the magnitude of the data keys of the first and second portions **112**, **118**, **132**, **138**, **152**, **158**. Other data values in the data elements that were not used to sort the data sets **110**, **130**, **150** may vary in magnitude between the values of first portions **112**, **132**, **152** and the second portions **118**, **138**, **158**.

By maintaining the magnitudes of the first and second portions **112**, **118**, **132**, **138**, **152**, **158**, the first and second portions **112**, **118**, **132**, **138**, **152**, **158** may be processed independently and/or in parallel. In a particular example, the first and second portions may be sorted in parallel (e.g., the processing time intervals of the portions may occur in substantially overlapping time intervals) and then merged to form a sorted data set that includes the data keys in the three data sets **110**, **130**, **150**. By maintaining the magnitude of the first and second portions **112**, **118**, **132**, **138**, **152**, **158** no additional sorting of the merged first and second portions may be necessary.

FIG. **2**A illustrates an example system **200** configured to process the partitioned portions **112**, **132**, **152**, **118**, **138**, **158** of the data sets **110**, **130**, **150** of FIG. **1**B, arranged in accordance with at least some embodiments described herein. The system **200** may include first and second memory units **210**, **212**. The first and second memory units **210**, **212** may be coupled to a first processor **220**. The second memory unit **212** may also be coupled to a second processor **222**.

FIG. **2**A illustrates the first portions **112**, **132**, **152** of the data sets **110**, **130**, **150** stored in the first memory unit **210** and the second portions **118**, **138**, **158** of the data sets **110**, **130**, **150** stored in the second memory unit **212**. In some embodiments, the first and second memory units **210**, **212** may be separate and physically distinct memory units. In other embodiments, the first and second memory units **210**, **212** may be separate logical partitions within the same memory unit.

In some embodiments, the first processor **220** may be configured to evaluate the data sets **110**, **130**, **150** to determine if the data sets **110**, **130**, **150** may be partitioned. For example, the first processor **220** may determine if the data sets **110**, **130**, **150** are sorted and that the data sets **110**, **130**, **150** each contain data elements with at least one data key that may be compared.

After evaluating the data sets **110**, **130**, **150**, the first processor **220** may be configured to identify partition indexes to partition the data sets **110**, **130**, **150** to form first portions **112**, **132**, **152** and second portions **118**, **138**, **158**. The first processor **220** may also be configured to load the first portions **112**, **132**, **152** and the second portions **118**, **138**, **158** into the first and second memory units **210**, **212**, respectively. The first processor **220** may also be configured to process the first portions **112**, **132**, **152** of the data sets **110**, **130**, **150**. The second processor **222** may be configured to process the second portions **118**, **138**, **158** of the data sets **110**, **130**, **150** upon receiving instructions from the first processor **220**. In some embodiments, the first portions **112**, **132**, **152** may be processed by the first processor **220** independently of the second portions **118**, **138**, **158**, which may be processed by the second processor **222**. In some embodiments, the first portions **112**, **132**, **152** may be processed in substantially overlapping time intervals with the second portions **118**, **138**, **158**. In some embodiments, the first and second processors **220**, **222** may be physically separate processors. In other embodiments, the first and second processors **220**, **222** may each be a different core in a multi-core processor.

In some embodiments, the first and second processors **220**, **222** may respectively be configured to merge and/or sort the first portions **112**, **132**, **152** and the second portions **118**, **138**, **158**. Any known sorting algorithm may be used. For example, the sorting algorithm may be a selection sort, insertion sort, comb sort, merge sort, heap sort, quick sort, counting sort, radix sort, or other type of sorting algorithm. In some embodiments, after sorting the second portions **118**, **138**, **158** the second processor **222** may merge the second portions **118**, **138**, **158** and indicate to the first processor **220** that the second portions **118**, **138**, **158** are sorted and merged. After receiving an indication that the second portions **118**, **138**, **158** are sorted and after sorting and merging the first portions **112**, **132**, **152**, the first processor **220** may combine the sorted and merged first portions **112**, **132**, **152** and second portions **118**, **138**, **158**. The first processor **220** may merge the sorted and merged first portions **112**, **132**, **152** and second portions **118**, **138**, **158** in either the first or second memory units **210**, **212** or in another memory unit.

By partitioning the data sets **110**, **130**, **150** into the first portions **112**, **132**, **152** and the second portions **118**, **138**, **158**; merging and sorting the first portions **112**, **132**, **152** and the second portions **118**, **138**, **158** in substantially overlapping time intervals; and then merging the merged and sorted first portions **112**, **132**, **152** and second portions **118**, **138**, **158**; the data sets **110**, **130**, **150** may be combined to form a sorted data set faster than might otherwise be possible. In particular, when the data sets **110**, **130**, **150** are large data sets, such as data sets that contain a number of data keys larger than about 50,000; 100,000; 500,000; 1,000,000; 20,000,000; or about 50,000,000; the reduced time to merge and sort the data sets **110**, **130**, **150** may be significant compared to merging and sorting the data sets **110**, **130**, **150** using a single processor or a single processor core.

FIG. **2**B illustrates another example system **250** configured to process the partitioned portions **112**, **132**, **152**, **118**, **138**, **158** of the data sets **110**, **130**, **150** of FIG. **1**B, arranged in accordance with at least some embodiments described herein. The system **250** includes a memory unit **260** that stores the first portions **112**, **132**, **152** and the second portions **118**, **138**, **158**. The system **250** further includes first, second, and third processors **270**, **272**, **274** networked together and to the memory unit **260**. The memory unit **260** and the processors **270**, **272**, **274** may be part of a computing cloud **280**.

In some embodiments, the third processor **274** may upload the first portions **112**, **132**, **152** and the second portions **118**, **138**, **158** into the memory unit **260** after they are sent to the computing cloud **280**. In some embodiments, the third processor **274** may upload the data sets **110**, **130**, **150** into the memory unit **260** and partition the data sets **110**, **130**, **150** to form the first portions **112**, **132**, **152** and the second portions **118**, **138**, **158**. In some embodiments, each processor **270**, **272**, **274** may have a separate memory unit. In these and other embodiments, the third processor **274** may upload the data sets **110**, **130**, **150** into each separate memory unit for partitioning by each memory unit's corresponding processor **270**, **272**, **274**.

In some embodiments, the third processor **274** may instruct the first processor **270** to process the first portions **112**, **132**, **152**. The third processor **274** may also instruct the second processor **272** to process the second portions **118**, **138**, **158**. The first and second processors **270**, **272** may process the first portions **112**, **132**, **152** and the second portions **118**, **138**, **158** respectively. In some embodiments, either of the first or second processors **270**, **272** may perform the functions of the third processor **274** and the third processor **274** may be omitted from the system **250**.

FIG. **1**A illustrates an example of 3-pair aligned data sets and provides various properties of the 3-pair aligned data sets. Generally, an aligned data set (ADS: Aligned Data Set) may be expressed based on the following notation ADS (name) =ADS(DDS:N). The left part of the notation defines a data set name and the right-hand side defines the composition of the data set. DDS is the abbreviation of Definition of Data Set, and defines the data elements within the data set in arbitrary

form. N specifies the number of data elements in the data set. 3 pair aligned data sets may be described as follows:

ADS3(name)=ADS3{ADS(DDS:*N*),ADS(DDS:*N*), ADS(DDS:*N*)}

The data sets may have the same number of data elements or a different number of data elements. The data elements within the 3-pair aligned data sets may each have a data key. The 3-pair aligned data sets may be organized in a logical order based on the data keys within the data elements of the 3-pair aligned data sets.

As noted before, the position of a data element within one of the data sets of the 3-pair aligned data sets may be specified using an index. In some embodiments, the index may commence at zero at the left hand side of the data set and increment by 1. In some embodiments, the index may commence at another number, such as 1, at either the left or right hand side, and may increment by 1 or some other value, such as, 2, 3, 4, 5, or some other value. Thus, the value of the index may indicate a data element within a data set. The data element may contain various data values and a data key.

In order to maintain a magnitude of data keys in the 3-pair aligned data sets when partitioning the 3-pair aligned data sets, locations to partition the 3-pair aligned data sets, referred to herein as partition indexes, may not be arbitrarily chosen. A method to partition the 3-pair aligned data sets may be derived based on a definition of the partition of the 3-pair aligned data sets. The method may commence by selecting an initial partition index and by comparing key values of the data elements at the initial partition index and indexes surrounding the initial partition index.

As an example, a first data set ADS (A:N), a second data set ADS (B:N), and a third data set ADS (C:N) may form 3-pair aligned data sets referred to as ADS3 {ADS (A:N), ADS (B:N), ADS (C:N)}, where A, B, and C respectively is an identifier of each of the respective data sets and N is the number of data elements in each of the data sets.

A proper partition of the 3-pair aligned data sets may result in 3 first portion aligned data sets (ADS3 (First Portion)) and 3 second portion aligned data sets (ADS3 (Second Portion)). When the 3-pair aligned data sets are properly partitioned:

$$\min[K\{\mathrm{ADS}(A(na1-1))\},K\{\mathrm{ADS}(B(nb1-1))\},K\{\mathrm{ADS}(C(nc1-1))\}]\geq\max[K\{\mathrm{ADS}(A(na1))\},K\{\mathrm{ADS}(B(nb1))\},K\{\mathrm{ADS}(C(nc1))\}] \text{ and}$$

$$na1+nb1+nc1=3x$$

where na1 is an arbitrary index value less than N−1 where the partition occurs in the A data set, nb1 is an arbitrary index value less than N−1 where the partition occurs in the B data set, and nc1 is an arbitrary index value less than N−1 where the partition occurs in the C data set with the indexing of the 3-pair aligned data sets commencing at zero. The variable 3x may be the number of data elements in the 3 first portions of the 3-pair aligned data sets where x is equal to (na1+nb1+nc1)/3.

Based on the above description of the partition, ADS3(First Portion)=ADS3 {ADS(A1:na1),ADS(B1:nb1),ADS(C1:nc1),3x} where A1 is an identifier for a first portion of the A data set, B1 is an identifier for a first portion of the B data set and C1 is an identifier for a first portion of the C data set. Furthermore, ADS3(Second Portion)=ADS3 {ADS(A2:N−na1),ADS(B2:N−nb1),ADS(C2:N−nc1),3N−3x} where A2 is an identifier for a second portion of the A data set, B2 is an identifier for a second portion of the B data set, and C2 is an identifier for a second portion of the C data set.

From the definitions, a first and a second formula may be derived as follows:

First formula—(na1−x)+(nb1−x)+(nc2−x)=0 and

Second formula—{(na1−1)−(x−1)}+{(nb1−1)−(x−1)}+{(nc2−1)−(x−1)}=0, where n(value1)(value2) represents a number of data elements, the value1 indicates the data set and the value2 indicates the portion of the partitioned data set. For example, nb1 represents the number of data elements in the first portion of the data set B.

The first and second formulas indicate that a combined difference between partition index values of tail data elements of the first portions of the 3-pair aligned data sets and an initial partition index value and a combined difference between partition index values of head data elements of the second portions of the 3-pair aligned data sets and the initial partition index value may be the same. Furthermore, a first property may be derived from the first and second formulas that the partition index the farthest from the initial partition index is offset from the initial partition index in a different direction than the remaining two partition indexes. For example, if the initial partition index is 10, and the partition index the farthest from the initial partition index is 6, then the remaining two partition indexes have values greater than or equal to 10.

A second property may be that a relationship between magnitudes of the data key of tail data elements of the first portions is unknown and a relationship between magnitudes of the data keys of head data elements of the second portions is unknown. The second property may be based on the distribution of the data key within each of the 3-pair aligned data sets being independent of each other.

A third property may be that when comparing the data key of data elements at the initial partition indexes and indexes surrounding the initial partition index and adjusting the indexes of a data set to compare the data key at the adjusted indexes, a first data set of the 3-pair aligned data sets with the largest data key and a second data set of the 3-pair aligned data sets with the smallest data key during one comparison may reverse during a next subsequent comparison so that the first data set may have the smallest data key and the second data set may have the largest data key.

FIGS. **3**A, **3**B, and **3**C show an example flow diagram of a method **300** of determining a partition index for each of three sorted data sets, arranged in accordance with at least some embodiments described herein. The method **300** may be performed in whole or in part on data sets, such as the data sets illustrated in FIG. **1**A and discussed above. The method **300** may be used to determine partition indexes where the data sets may be partitioned. The method **300** may include various operations, functions, or actions as illustrated by one or more of blocks **302**, **304**, **306**, **308**, **310**, **312**, **314**, **316**, **318**, **320**, **330**, **332**, **334**, **336**, **338**, **340**, **342**, **350**, **352**, **354**, **356**, **358**, **360**, and/or **362** in FIGS. **3**A, **3**B, and **3**B. Referring now to FIG. **3**A, the method **300** may begin at block **302**.

In block **302**, ["Select Initial Partition Indexes"], a computing device may be configured (e.g., via software or firmware) to select initial partition indexes for partitioning the data sets. The data sets may include data set A, data set B, and data set C. In some embodiments, the initial partition indexes may be the same for each data set. In some embodiments, the initial partition indexes may be different for each data set. In some embodiments, the initial partition indexes may be determined based on the number of data elements within each of the data sets. For example, the initial partition indexes may be equal to the modulus of the combined number of data elements in all three sets divided by 2, 3, 4, 5, 6 or some other number.

The location of the initial partition indexes may also vary based on the number of portions into which the data sets are

partitioned and based on approximately how many data elements are desired in each portion of a data set. For example, in some embodiments, it may be desired to have the number of data elements in each portion be equal or substantially equal. In some embodiments, the number of data elements in the first portions based on the initial partition indexes may be a multiple of 3. In these and other embodiments, the initial partition index may be x−1 for each data set, where 3x equals the number of data elements in the first portions. In some embodiments, the number of data elements in the first portions based on the initial partition indexes may not be a multiple of three. In these and other embodiments, the number of data elements may be 3x+1 or 3x+2. In the case of 3x+1 number of data elements, the initial partition indexes are selected as x. The partition index of the data set with the largest data key at the initial partition index of x remains x. The other initial partition indexes for the other data sets may be x−1. In the case of 3x+2 number of data elements, the initial partition indexes are selected as x. The partition index of the data set with the smallest data key at the initial partition index of x remains x. The other initial partition indexes for the other data sets may be x−1. In some embodiments, the initial partition indexes may vary based on whether head data elements of the second portions of the data sets or tail data elements of the first portions of the data sets are being determined. For example, when determining the tail data elements, the initial partition indexes may be x−1 and when determining the head data elements the initial partition indexes may be x.

The initial partition indexes of the data sets may be represented by (i, j, k) with i being the partition index for data set A, j being the partition index for data set B, and k being the partition index for data set C. The key value of the data set A at the partition index i may be represented by A(i). The key value of the data set B at the partition index j may be represented by B(j). The key value of the data set C at the partition index k may be represented by C(k). Block **302** may be followed by block **304**.

In block **304**, ["Are Data Keys at the Initial Partition Indexes Equal?"], the computing device can be configured to compare the data keys at the partition indexes of i, j, k, which are equal to the initial partition indexes set in block **302**, to determine if the data keys are all substantially equal. If the data keys are determined to be substantially equal to one another, then block **304** may be followed by block **306**. If the data values are determined to not be substantially equal to one another, then block **304** may be followed by block **308**.

In block **306**, ["Final Partition Indexes Determined"], the computing device may be configured to determine that the final partition indexes for the data sets may equal the initial partition indexes used for the data sets. The method **300** may be complete.

In block **308**, ["Compare Data Keys At The Initial Partition Indexes"], the computing device may be configured to determine the data set with the largest data key, the data set with the smallest data key, and the data set with the middle data key based on the initial partition indexes. Block **308** may be followed by block **310**.

In block **310**, ["Set Old Comparison Code"], the computing device may be configured to set an old comparison code (C_old) based on key values of the various data sets at the initial partition indexes. For example, when A(i)>B(j)>C(k) then C_old may equal 1. When A(i)>C(k)>B(j) then C_old may equal 2. When B(j)>A(i)>C(k) then C_old may equal 3. When B(j)>C(k)>A(i) then C_old may equal 4. When C(k)>A(i)>B(j) then C_old may equal 5. When C(k)>B(j)>A(i) then C_old may equal 6. The C_old may also contain the index values for the data set A, the data set B, and the data set C that are used to set the C_old.

In block **312**, ["Update Partition Indexes"] the computing device may be configured to adjust (e.g., increment or decrement) the partition indexes for the data sets with the highest and lowest data keys at their respective partition indexes to generate updated partition indexes. For example, when A(i)>B(j)>C(k), then the partition index of the data set A and the data set C may be updated by adjusting the partition indexes i and k. In these and other embodiments, the partition index j of data set B may not be adjusted.

In some embodiments, the partition indexes may be adjusted by incrementing or decrementing (e.g., by one, two or some other increment value) to an adjacent or nearby index. The partition index of the data set with the highest data key may be adjusted to a partition index with a lower data key. Similarly, the partition index for the data set with the lowest data key may be adjusted to a partition index with a higher data key. When the partition index for the data set with the highest data key is incremented, the partition index for the data set with the lowest data key may be decremented. Alternately, when the partition index for the data set with the highest data key is decremented, the partition index for the data set with the lowest data key may be incremented.

In some embodiments, when the partition index is adjusted (e.g., incremented or decremented) the adjusted partition index may fall outside the number of data keys in a data set. If the current partition index of a data set is adjusted beyond the number of data keys in the data set, a random data value may be associated with the data set so that future comparisons of the value of the particular data set at the adjusted partition index may occur. Associating the random value with the particular data set does not insert the random value into the data set. Rather, the random value is merely associated with the data set for comparison purposes within the method **300**. For example, when the data key at a partition index is the lowest data key just before adjusting the partition index beyond the number of data elements in the data set, the random data key associated with the data set may be lower than any data key in any of the data sets. Conversely, when the data key at a partition index is the largest data key in a data set just before adjusting the partition index beyond the number of data keys in the data set, the random data key associated with the data set may be larger than any data key in any of the data sets. Block **312** may be followed by block **314**.

In block **314**, ["Compare Data Keys At The Updated Partition Indexes"], the computing device may be configured to determine the data set with the largest data key, the data set with the smallest data key, and the data set with the middle data key based on the updated partition indexes of the data sets. Block **314** may be followed by block **316**.

In block **316**, ["Set New Comparison Code"], the computing device may be configured to set a new comparison code (C_new) based on key values of the various data sets at the updated partition indexes. For example, when A(i)>B(j)>C(k) then C_new may equal 1. When A(i)>C(k)>B(j) then C_new may equal 2. When B(j)>A(i)>C(k) then C_new may equal 3. When B(j)>C(k)>A(i) then C_new may equal 4. When C(k)>A(i)>B(j) then C_new may equal 5. When C(k)>B(j)>A(i) then C_new may equal 6. The C_new may also contain the partition index values for the data set A, the data set B, and the data set C that are used to set the C_new. Block **316** may be followed by block **318**.

In block **318**, ["Reversal Occurred?"], the computing device may be configured to determine if a reversal occurred based on the definition of the third property, namely that a first data set of the 3-pair aligned data sets with the largest key

value and a second data set 3-pair aligned data sets with the smallest key value during one comparison are reversed so that the first data set has the smallest key value and the second data set has the largest key value during a next subsequent comparison. The reversal may be determined by comparing the C_old and the C_new. A reversal occurs when C_new equals 1 and C_old equals 6, when C_new equals 2 and C_old equals 4, when C_new equals 3 and C_old equals 5, when C_new equals 4 and C_old equals 2, when C_new equals 5 and C_old equals 3, when C_new equals 6 and C_old equals 1. As an example, when C_old equals 6, C(k)>B(j)>A(i) at a previous partition index. When C_new equals 1, A(i)>B(j)>C(k) at an updated partition index. A reversal occurs because C(k) is the greatest at the previous partition index and the least at the updated partition index and A(i) is the least at the previous partition index and the greatest at the updated partition index.

When a reversal occurs, the method **300** may be followed by block **330** (see FIG. **3**B) for selection of final partition indexes for the tail data elements of the first portions of the data sets and/or by block **350** (see FIG. **3**C) for selection of final partition indexes for the head data elements of the second portions of the data sets. When a reversal does not occur, the method **300** may be followed by block **320**.

In block **320**, ["Set New Comparison Code to Old Comparison Code"], the computing device may be configured to set the C_new to the C_old, such that the C_old contains the value of the C_new. For example, if C_new equals 4 and C_old equals 1, after setting C_new to C_old, C_new and C_old would both equal 4. Block **320** may be followed by block **312**. Blocks **312**, **314**, **316**, **318**, and **320** may be repeated until a reversal occurs in block **318**.

Referring now to FIG. **3**B, in block **330**, ["Determine Tail Candidate **1**"], the computing device may be configured to determine a partition index of a data set for tail candidate **1**. The tail candidate **1** may be the partition index of the data set with the highest data key of the partition indexes of the data sets with the lowest data keys from the C_old and the C_new that determined a reversal occurred in block **318**. For example, when the C_old equals 6 based on C(k)>B(j)>A(i) at a previous partition index and the C_new equals 1 based on A(i)>B(j)>C(k) at an updated partition index, the A(i) at the previous partition index is compared to C(k) at the updated partition index to determine which is larger. The partition index with the larger data key is selected as the tail candidate **1**. Block **330** may be followed by block **332**.

In block **332**, ["Determine Tail Candidate 2"], the computing device may be configured to determine a partition index of a data set for tail candidate **2**. Tail candidate 2 may be the partition index adjacent to a partition index of a data set that has a middle data key from the C_old and the C_new and that has a data key less than the middle data key. For example, assuming the C_old and the C_new equals 6 and 1 respectively, the partition index of the data set that is the middle data key is j of B(j) and the tail candidate **2** is the partition index adjacent to j that has a data key less than B(j). Block **332** may be followed by block **334**.

In block **334**, ["Compare Data Key At Tail Candidate 1 With Data Key at Tail Candidate 2"], the computing device may be configured to compare the data key at the tail candidate **1** with the data key at the tail candidate **2**. When the data key at the tail candidate **1** is larger, block **334** may be followed by block **336**. When the data key at the tail candidate **2** is larger, block **334** may be followed by block **340**.

In block **336**, ["Assign Tail Candidate **1** To First Final Partition Index"], the computing device may be configured to assign the tail candidate **1** to a first final partition index. Block **336** may be followed by block **338**.

In block **338**, ["Determine Second And Third Final Partition Indexes"], the computing device may be configured to determine a second and a third final partition indexes. The second final partition index may be the partition index of the data set that is the middle data key from the C_old and the C_new that resulted in a reversal in block **318**. The third final partition index may be the partition index adjacent to the partition index of the data set that has the lowest data key from the C_old and the C_new and that has a data key larger than the lowest data key. For example, when the C_old equals 6 based on C(k)>B(j)>A(i) at a previous partition index and the C_new equals 1 based on A(i)>B(j)>C(k) at an updated partition index, the A(i) at the previous partition index is compared to C(k) at the updated partition index to determine which is smaller. The partition index adjacent to the partition index with the smaller data key that has a data key larger than the smaller data key is selected as the third final partition index.

In block **340**, ["Assign Tail Candidate 2 To First Final Partition Index"], the computing device may be configured to assign the tail candidate **2** to a first final partition index. Block **340** may be followed by block **342**.

In block **342**, ["Determine Second And Third Final Partition Indexes"], the computing device may be configured to determine a second and a third final partition indexes. The second final partition index may be the partition index adjacent to the tail candidate **1** that has a data key larger than the data key at the tail candidate **1**. The third final partition index may be the partition index adjacent to the partition index of the data set that has the lowest data key from the C_old and the C_new and that has a data key larger than the lowest data key. For example, when the C_old equals 6 based on C(k)>B(j)>A(i) at a previous partition index and the C_new equals 1 based on A(i)>B(j)>C(k) at a current index, the A(i) at a previous partition index is compared to C(k) at an updated index to determine which is smaller. The partition index adjacent to the partition index with the smaller data key that has a data key larger than the smaller data key is selected as the third final partition index.

Referring now to FIG. **3**C, in block **350**, ["Determine Head Candidate **1**"], the computing device may be configured to determine a partition index of a data set for head candidate **1**. The head candidate **1** may be the partition index of the data set with the lowest data key of the partition indexes of the data sets with the highest data keys from the C_old and the C_new. For example, when the C_old equals 6 based on C(k)>B(j)>A(i) at a previous partition index and the C_new equals 1 based on A(i)>B(j)>C(k) at an updated partition index, the C(k) at the previous partition index is compared to A(i) at the updated partition index to determine which is smaller. The partition index with the smaller data key is selected as the head candidate **1**. Block **350** may be followed by block **352**.

In block **352**, ["Determine Head Candidate **2**"], the computing device may be configured to determine a partition index of a data set for head candidate **2**. Head candidate **2** may be the partition index adjacent to a partition index of a data set that is a middle data key from the C_old and the C_new and that has a data key greater than the middle data key. For example, assuming the C_old and the C_new equal 6 and 1 respectively, the partition index of the data set that is the middle data key is j of B(j) and the head candidate **2** is the partition index adjacent to j that has a data key greater than B(j). Block **352** may be followed by block **354**.

In block **354**, ["Compare Data Key at Head Candidate **1** With Data Key at Head Candidate **2**"], the computing device may be configured to compare the data key at head candidate **1** with the data key of head candidate **2**. When the data key at

head candidate **1** is larger, block **354** may be followed by block **356**. When the data key at head candidate **2** is larger, block **354** may be followed by block **360**.

In block **356**, ["Assign Head Candidate **2** To First Final Partition Index"], the computing device may be configured to assign the head candidate **2** to a first final partition index. Block **356** may be followed by block **358**.

In block **358**, ["Determine Second And Third Final Partition Indexes"], the computing device may be configured to determine a second and a third final partition indexes. The second final partition index may be the partition index with the smallest data key that is compared with the head candidate **1** when the head candidate **1** has the largest data key from the C_old or the C_new. For example, when the C_old is C(k)>B(j)>A(i) and the head candidate **1** is k of C(k) then the second final partition index is i of A(i). The third final partition index may be a partition index adjacent to the head candidate **1** that has a data key less than the data key at the head candidate **1**.

In block **360**, ["Assign Head Candidate 1 To First Final Partition Index"], the computing device may be configured to assign the head candidate **1** to a first final partition index. Block **360** may be followed by block **362**.

In block **362**, ["Determine Second And Third Final Partition Indexes"], the computing device may be configured to determine a second and a third final partition index. The second final partition index may be the partition index with the smallest data key that is compared with partition index of the head candidate **1** when the partition index of the head candidate **1** has the largest data key from the C_old or the C_new. For example, when the C_old is C(k)>B(j)>A(i) and head candidate **1** is k of C(k) then the second final partition index is i of A(i). The third final partition index may be a partition index adjacent to the head candidate **2** that has a data key less than the data key at the head candidate **2**.

As described, the method **300** may be used to determine a partition index for each of three sorted data sets. The partition indexes may be used by a computing device to partition each of three sorted data sets into first and second portions where each data value in each of the first portions of the corresponding data sets is greater in magnitude than each data value in each of the second portions of the corresponding data sets. In some embodiments, the computing device may partition each of three sorted data sets using a single processing thread. In these and other embodiments, the processing thread may be passed information concerning the three sorted data sets, such as a pointer to a beginning or ending index of each of the three sorted data sets. Additional information may also be passed to the processing thread, such as an initial partition index.

In some embodiments, the data elements in the data sets may have multiple data keys. In these and other embodiments, a first data key of the data elements may be used for comparison within the method **300**. When the first data key of data elements are equal, additional comparisons of the additional data keys may be made within the method **300**.

In other embodiments, the method **300** may be implemented on three sorted data sets to partition each of the three-sorted data sets into three portions. More generally, the described method **300** may be implemented on three sorted data sets to partition each of the three-sorted data sets into any number of portions. In these embodiments and others where the sorted data sets may be partitioned into more than two portions, the method **300** may be performed by a single processor multiple times to partition the data sets into portions. Alternately or additionally, the method **300** may be performed by more than one processor independently and/or in parallel. For example, if each of the sorted data sets are partitioned into three portions, a first processor may be configured to determine the partition indexes to form the first and second portions and a second processor may be configured to determine the partition indexes to form the third portions.

In some embodiments, determining the partition indexes to form the first and second portions and to form the third portions may be performed by the first and second processors, respectively, in parallel and in substantially overlapping time intervals. As another example, if each of the sorted data sets are partitioned into 100 portions, 99 different processors may be configured to determine the 99 partition indexes to form the 100 portions. Thus, the method **300** may be performed using numerous processor, such as 100, 200, 500, 1000, or more processors. The method **300** may be performed independently and/or in parallel by multiple processors because the method **300** involves reading and comparing the data values at the different indexes of the sorted data sets and not changing, alternating, or adjusting data values within the sorted data sets.

Some embodiments disclosed herein may include a computer storage medium having computer-executable instructions stored thereon that are executable by a computing device to perform operations including but not limited to those operations described with respect to method **300** of FIGS. **3**A-**3**C, such as the operations illustrated by blocks **302**, **304**, **306**, **308**, **310**, **312**, **314**, **316**, **318**, **320**, **330**, **332**, **334**, **336**, **338**, **340**, **342**, **350**, **352**, **354**, **356**, **358**, **360**, **362**, or some combination thereof in FIGS. **3**A, **3**B, and **3**C, and/or variations thereof.

One skilled in the art will appreciate that, for this and other processes and methods disclosed herein, the functions performed in the processes and methods may be implemented in differing order. Furthermore, the outlined steps and operations are only provided as examples, and some of the steps and operations may be optional, combined into fewer steps and operations, or expanded into additional steps and operations without detracting from the essence of the disclosed embodiments.

For example, in some embodiments blocks **350**-**362** may not be performed to determine the heads of the second portions of the data sets. In some embodiments, blocks **330**-**342** may not be performed to determine the tails of the first portions of the data sets.

FIG. **4**A illustrates three example data sets, data set A, data set B, and data set C, arranged in accordance with at least some embodiments described herein. Data set A contains 6 data elements arranged in descending order and is associated with an index value i. Data set B contains 6 data elements arranged in descending order and is associated with an index value j. Data set C contains 6 data elements arranged in descending order and is associated with an index value k. An example of the method **300** illustrated in FIGS. **3**A-**3**C is now described with reference to the data set A, the data set B, and the data set C of FIG. **4**A.

In block **302**, the initial partition indexes are selected based on the number of data elements in the data sets. The initial partition indexes are selected as 1, where x=2 and 3x=6. Thus, i=j=k=1 where i is the partition index of the data set A, j is the partition index of the data set B, and k is the partition index of the data set C.

In block **304**, the data keys at the initial partition indexes, A(1)=290, B(1)=260, and C(1)=265, are determined to not be equal, such that block **304** is followed by block **308**. In block **308**, the data keys from the data sets are compared based on the initial partition indexes and it is determined that A(1)=290>C(1)=265>B(1)=260. Block **308** is followed by block **310**. In block **310**, the C_old is set to 2. Block **310** is followed by block **312**.

In block **312**, the partition index i of data set A and the partition index j of the data set B is adjusted because data set A has the highest data key at a current partition index and data set B has the lowest data key at the current partition index. The partition indexes are adjusted so that i=2, j=0. The partition index of k is not adjusted because the data set C has the middle data key at the current partition index. Block **312** is followed by block **314**. In block **314**, the data keys from the data sets are compared based on the adjusted partition indexes and it is determined that B(0)=280>C(1)=265>A(2)=262. Block **314** is followed by block **316**. In block **316**, the C_new is set to 4. Block **316** is followed by block **318**. In block **318**, it is determined that a reversal occurred because the C_new equals 4 and the C_old equals 2. Block **318** may be followed by blocks **330** and/or **350**. In this example, block **318** is followed by block **330**.

In block **330**, the partition index with the smallest data key (260) from the C_old is j=1 (B(1)=260) and the partition index with the smallest data key (262) from the C_new is i=2 (A(2)=262). Since the data key at i=2 of the data set A is larger than the data key at j=1 of the data set B, i=2 is thus determined to be the first tail candidate. Block **330** may be followed by block **332**. In block **332**, the partition index of the data set with the middle data key (265) in the C_old and the C_new is k=1. The partition index adjacent to k=1 that contains a data key less than the data key at k=1 is k=2, so k=2 is determined to be the second tail candidate. Block **332** may be followed by block **334**. In block **334**, the data key (262) at the tail candidate **1** of i=2 is determined to be greater than the data key (245) at the tail candidate **2** of k=2. Thus, block **334** may be followed by block **336**.

In block **336**, the first tail candidate i=2 is determined as the first final partition index and as the partition index of the data set A. Thus, the data element that contains the data key of 262 is the tail data element of the first portion of the data set A. Block **336** may be followed by block **338**. In block **338**, the second final partition index may be k=1 because k=1 is the partition index of the data set with the middle data key in the C_old and the C_new. Thus, the data element that contains the data key of 265 is the tail data element of the first portion of the data set C. The partition index with the lowest data keys from the C_old and the C_new are determined to be i=2 and j=1. Of i=2 and j=1, j=1 has a lower data key. The partition index adjacent to j=1 that has a higher data key than j=1 is j=0. Thus, j=0 is the third final partition index and the data element that contains the data key of 280 is the tail data element of the first portion of the data set B. FIG. **4**B illustrates an example partitioning of the three example data sets of FIG. **4**A using the method **300**, arranged in accordance with at least some embodiments described herein.

FIG. **5**A illustrates three example data sets, data set A, data set B, and data set C, arranged in accordance with at least some embodiments described herein. Data set A contains 6 data elements arranged in descending order and is associated with an index value i. Data set B contains 6 data elements arranged in descending order and is associated with an index value j. Data set C contains 6 data elements arranged in descending order and is associated with an index value k. An example of the method **300** illustrated in FIGS. **3**A-**3**C is now described with reference to the data set A, the data set B, and the data set C of FIG. **5**A.

In block **302**, the initial partition indexes are selected based on the number of data elements in the data sets. The initial partition indexes are selected as 2. Thus, i=j=k=2, where i is the partition index of data set A, j is the partition index of data set B, and k is the partition index of data set C.

In block **304**, the data keys at the initial partition indexes, A(2)=180, B(2)=172, and C(2)=135, are determined to not be equal, such that block **304** is followed by block **308**. In block **308**, the data keys from the data sets are compared based on the initial partition indexes and it is determined that A(2)=180>B(2)=172>C(2)=135. Block **308** is followed by block **310**. In block **310**, the C_old is set to 1. Block **310** is followed by block **312**.

In block **312**, the partition indexes of the data set A and the data set C are adjusted because the data set A has the highest data key (180) at a current partition index and the data set C has the lowest data key (135) at the current partition index. The partition indexes are adjusted so that i=3, k=1. The partition index of j is not adjusted because the data set B has the middle data key (172) at the current partition index. Block **312** is followed by block **314**. In block **314**, the data keys from the data sets are compared based on the adjusted partition indexes and it is determined that B(2)=172>A(3)=170>C(1)=145. Block **314** is followed by block **316**. In block **316**, the C_new is set to 3. Block **316** is followed by block **318**. In block **318**, it is determined that a reversal did not occur because the C_new equals 3 and the C_old equals 1. Block **318** may be followed by block **320**. In block **320**, the C_new is set to the C_old so that the C_old equals 3. Block **320** may be followed by block **312**.

In block **312**, the partition indexes of the data set B and the data set C are adjusted because the data set B has the highest data value (**172**) at the current partition index and the data set C has the lowest data value (**145**) at the current partition index. The partition indexes are adjusted so that j=3, k=0. The partition index of i is not adjusted because the data set A has the middle data key (170) at the current partition index. Block **312** is followed by block **314**. In block **314**, the data keys from the data sets are compared based on the adjusted partition indexes and it is determined that B(3)=171>A(3)=170>C(0)=155. Block **314** is followed by block **316**. In block **316**, the C_new is set to 3. Block **316** is followed by block **318**. In block **318**, it is determined that a reversal did not occur because the C_new equals 3 and the C_old equals 3. Block **318** may be followed by block **320**. In block **320**, the C_new is set to the C_old so that the C_old equals 3. Block **320** may be followed by block **312**.

In block **312**, the partition indexes of the data set B and the data set C are adjusted because the data set B has the highest data value (**171**) at the current partition index and the data set C has the lowest data value (**155**) at the current partition index. The partition index of the data set B is adjusted so that j=4. Adjusting the partition index of the data set C would cause the partition index to fall outside the number of data keys in a data set. Thus, a random data key of +infinity may be associated with the data set for future comparisons of the data set C. The index of the random data key may be assigned −1 for exemplary purposes The partition index of i is not adjusted because data set A has the middle data key (170) at the current partition index. Block **312** is followed by block **314**. In block **314**, the data keys from the data sets are compared based on the adjusted partition indexes and it is determined that C(−1)=+inf>A(3)=170>B(4)=140. Block **314** is followed by block **316**. In block **316**, the C_new is set to 5. Block **316** is followed by block **318**. In block **318**, it is determined that a reversal occurred because the C_new equals 5 and the C_old equals 3. Block **318** may be followed by blocks **330** and **350**. In this example, block **318** is followed by block **330** only.

In block **330**, the partition index with the smallest data key (155) from the C_old is k=0 (C(0)=155) and the partition index with the smallest data key (160) from the C_new is j=4 (B(4)=140). Since the data key at k=0 of the data set C is

larger than the data key at j=4 of the data set B, k=0 is thus determined to be the first tail candidate. Block **330** may be followed by block **332**. In block **332**, the partition index of the data set with the middle data key (170) in the C_old and the C_new is i=3. The partition index adjacent to i=3 that contains a data key less than the data key at i=3 is i=4, so i=4 is determined to be the second tail candidate. Block **332** may be followed by block **334**. In block **334**, the data key (155) at the tail candidate **1** of k=0 is determined to be smaller than the data key (160) at the tail candidate **2** of i=4. Thus, block **334** may be followed by block **340**.

In block **340**, the second tail candidate i=4 is determined as the first final partition index and as the partition index of the data set A. Thus, the data element that contains the data key of 160 is the tail data element of the first portion of the data set A. Block **340** may be followed by block **342**. In block **342**, the second final partition index may not exist because there is no partition index adjacent to the partition index of C(0) (tail candidate **1**), that has a data key larger than the data key of C(0). The partition index with the lowest data keys from the C_old and the C_new are determined to be k=0 and j=4. Of k=0 and j=4, j=4 has a lower data key. The partition index adjacent to j=4 that has a higher data key than j=4 is j=3. Thus, j=3 is the third final partition index and the data element that contains the data key of 171 is the tail data element of the first portion of the data set B. FIG. **5**B illustrates an example partitioning of the three example data sets of FIG. **5**A using the method **300**, arranged in accordance with at least some embodiments described herein.

FIG. **6** illustrates an example system **600** for implementing the method **300** of FIGS. **3**A-**3**C, arranged in accordance with at least some embodiments described herein. The system **600** may include an initial comparison unit **610**, a comparing unit **620**, a final partition index determination unit **630**, and various memory units including an index variables unit **612**, a C_old unit **614**, and a C_new unit **616**.

The system **600** may be configured to partition 3-pair aligned data sets (referred to herein as "data sets"), such as the data sets discussed with respect to FIGS. **1** and **3**A-**3**C, into first and second portions. The initial comparison unit **610** may be configured to determine an initial partition index for each of the data sets. In some embodiments, the initial partition index may be the same for each data set or one or more of the data sets may have a different initial partition index. In some embodiments, the initial partition indexes may vary based on the number of data elements in the data sets and/or whether head data elements and/or tail data elements for the first and/or second portions respectively are being determined. After determining the initial partition indexes, the initial comparison unit **610** may send the initial partition indexes to the index variables unit **612**.

The initial comparison unit **610** may further be configured to compare the data keys at the initial partition indexes to determine if the data keys are equal. When the data keys are equal, the initial comparison unit **610** may send the index variables to the final partition index determination unit **630**. When the data keys are not equal, the initial comparison unit **610** may determine which of the data keys is the largest, the smallest, and the middle and send the information along with partition index variables that are associated with the data keys to the C_old unit **614**. Further, when the data keys are not equal, the initial comparison unit **610** may adjust the partition indexes of the data sets with the largest and smallest data keys and send the adjusted partition indexes to the index variables unit **612**. The initial comparison unit **610** may indicate to the comparing unit **620** to begin.

The comparing unit **620** may be configured to access the partition indexes from the index variables unit **612** and compare data keys from each of the data sets at the accessed partition indexes. The comparing unit **620** may determine which of the data keys is the largest, the smallest, and the middle and send the information along with index variables that are associated with the data keys to the C_new unit **616**.

The comparing unit **620** may also be configured to determine when the data set with the largest key value in the C_old unit **614** is the data set with the smallest key value in the C_new unit **616** and when the data set with the smallest key value in the C_old unit **614** is the data set with the largest key value in the C_new unit **616**. This may be referred to as a reversal.

When the comparing unit **620** determines that a reversal has occurred, the comparing unit **620** may be configured to so indicate to the final partition index determination unit **630**. The comparing unit **620** may send to the final partition index determination unit **630** the information from the C_old unit **614** and the C_new unit **616** or some information derived therefrom. For example, in some embodiments, the comparing unit **620** may determine a reversal code indicating the type of reversal that occurred. For example, the reversal code may indicate which of the data sets changed data keys and which one of the data sets did not change its data key.

When the comparing unit **620** determines that a reversal has not occurred, the comparing unit **620** may be configured to set the C_old **614** equal to the C_new **616** and to adjust the partition indexes of the data sets with the largest and smallest data keys and send the adjusted partition indexes to the index variables unit **612**. The comparing unit **620** may then repeat by accessing the partition indexes from the index variables unit **612**, comparing the data keys from each of the data sets at the partition indexes, sending the information to the C_new unit **616**, and determining if a reversal occurs.

The final partition index determination unit **630** may be configured to receive the information from the comparing unit **620**, such as the reversal code, the information from the C_old unit **614** and the C_new unit **616**, and/or other information. The final partition index determination unit **630** may be configured to use the information to determine tail candidates **1** and **2** and/or head candidates **1** and **2**. Additionally, the final partition index determination unit **630** may be configured to determine the final partition indexes for the tail data elements and/or the head data elements based on the tail candidates **1** and **2** and/or head candidates **1** and **2**. In some embodiments, when the final partition index determination unit **630** receives a reversal code from the comparing unit **620**, the final partition index determination unit **630** may have different sections that may handle each of the different possible reversal codes.

The final partition index determination unit **630** may also be configured to receive the indication that the data keys are all equal from the initial comparison unit **610**. Based on this indication, the final partition index determination unit **630** may determine that the final partition indexes are equal to the initial partition indexes.

FIG. **7** illustrates an example merge **700** of three example data sets, arranged in accordance with at least some embodiments described herein. The merge **700** occurs of three data sets **710**, **730**, **750** that are combined into a merged data set **780**. The data sets **710**, **730**, **750** may be similar to the data sets **110**, **130**, **150** of FIG. **1**A. The data sets **710**, **730**, **750** may each be aligned data sets that include various sorted data elements and that together may form a 3 pair aligned data set. The data sets **710**, **730**, **750** may each be divided into nine portions. Indexes **770**-**777** may be used as initial partition

indexes to divide each of the data sets **710**, **730**, **750** into nine portions following the method described herein and illustrated in FIGS. **3**A-**3**C. The index **770** may be equal to N/9 where N is the number of data elements in at least one of the data sets **710**, **730**, **750**. The indexes **771**-**777** may be equal to 2N/9, 3N/9, 4N/9, 5N/9, 6N/9, 7N/9, and 8N/9, respectively.

Using the method described herein and illustrated in FIGS. **3**A-**3**C, the data set **710** may be divided into portions **711**-**719**, the data set **730** may be divided into portions **731**-**739**, and the data set **750** may be divided into portions **751**-**759**.

The portions **711**, **731**, and **751** may be sorted and combined to form a portion **781** of the merged data set **780**. Likewise, portions **712**, **732**, **752** may be sorted and combined to form a portion **782** of the merged data set **780**. Portions **713**, **733**, **753** may be sorted and combined to form a portion **783** of the merged data set **780**. Portions **714**, **734**, **754** may be sorted and combined to form a portion **784** of the merged data set **780**. Portions **715**, **735**, **755** may be sorted and combined to form a portion **785** of the merged data set **780**. Portions **716**, **736**, **756** may be sorted and combined to form a portion **786** of the merged data set **780**. Portions **717**, **737**, **757** may be sorted and combined to form a portion **787** of the merged data set **780**. Portions **718**, **738**, **758** may be sorted and combined to form a portion **788** of the merged data set **780**. Portions **719**, **739**, **759** may be sorted and combined to form a portion **789** of the merged data set **780**. The portions **781**-**789** are combined to form the merged data set **780**. The merged data set **780** is a sorted combination of the three data sets **710**, **730**, **750**. The combination of the data sets **710**, **730**, **750** using the method disclosed herein, where the time to divide the data sets **710**, **730**, **750** is disregarded, may take ⅑ of the time to form the merged data set **780** as compared to combining the data sets **710**, **730**, **750** using conventional methods, such as a quick sort method.

Although the data sets **710**, **730**, **750** are illustrated as being divided into 9 portions, depending on a number of data elements within the data sets **710**, **730**, **750**, the data sets **710**, **730**, **750** may be divided into more or less portions. For example, the data sets **710**, **730**, **750** may be divided into 50, 100, or 1000 portions.

FIG. **8** illustrates an example merge **800** of nine data sets **810**, arranged in accordance with at least some embodiments described herein. The nine data sets **810** may be formed from a non-sorted data set or may be nine separate data sets. The total number of data elements in the nine data sets **810** may be N. In some embodiments, each data set of the data sets **810** may have N/9 data elements. The nine data sets **810** may be sorted using a quick sort module **820**. After sorting, the nine data sets **810** may form nine sorted data sets **830**.

The nine sorted data sets **830** may be grouped together into 3 sets of 3-pair type aligned data sets. Two partitions indexes may be selected to divide the nine sorted data sets **830** each into three portions. The partition indexes may be (⅓)(N/9) and (⅔)(N/9). During a first stage, the portions of each set of the 3 pair type aligned data sets may be merged together as illustrated and described with respect to FIG. **7** to form data sets **840**, **842**, **844**. For example, three sorted data sets of the nine sorted data sets **830** that form one of the 3 pair type aligned data sets may be partitioned, each of the corresponding portions from the three sorted data sets may be sorted and merged, and the merged portions may be merged to form one of the data sets **840**, **842**, or **844**. Each of the merged portions may include N/9 data elements. In some embodiments, nine sets of threads may be used to perform a merge of three data sets in one set of 3 pair type aligned data sets. Six sets of threads may perform the sorting and merging after determining partition indexes for the three data sets. Three sets of threads may perform sorting and merging from the head of each of the three data sets during and after determining the partition indexes for the three data sets.

The data sets **840**, **842**, **844** may form a set of 3 pair aligned data sets. During a second stage, the data sets **840**, **842**, **844** may be merged together as illustrated and described with respect to FIG. **7** to form data set **870**. The partition indexes for partitioning the data sets **840**, **842**, **844** into nine portions may be (⅑)(N/3), (⅔)(N/3), (⅓)(N/3), (⅘)(N/3), (⅕)(N/3), (⅚)(N/3), (⅞)(N/3), (⅛)(N/3). Each of the merged portions may include N/9 data elements. In some embodiments, nine sets of threads may be used to perform a merge of data sets **840**, **842**, **844** into the data set **870**. Eight sets of threads may perform the sorting and merging after determining partition indexes for the three data sets. One set of threads may perform sorting and merging from the head of each of the data sets **840**, **842**, **844** during and after determining the partition indexes for the data sets **840**, **842**, **844**.

The data set **870** may be a sorted data set that includes the data elements from the nine data sets **810**. To combine the nine data sets **810** into one sorted data set using conventional methods, such as a quick sort method, may have a processing time of Tqs(N), where N is the number of data elements in the nine data sets **810**. The order of the processing time for the quick sort method may be $n \log_{2N}$. Thus, the processing time for a quick sort method increases in proportion to the number of elements being sorted.

The processing time for forming the data set **870** using the above described method may be much shorter. As described, the nine data sets **810** may be sorted in parallel and then merged together. Thus, the processing time may be equal to the time for sorting the nine data sets **810** and to perform the first and second stage of the merging as described above. The processing time to sort each of the nine data sets **810** in parallel may be equal to (⅑) the time required to sort the all of the data elements in the nine data sets **810** or Tqs(N/9). The processing time for performing the first stage of the merge may be equal to Tm(N/3). The processing time to perform the second stage of the merge may be equal to Tm(N). The processing times of Tm(N/3) and Tm(N) may be much shorter than the processing time for a quick sort because the position of a data element may be decided using two comparison operations where the position of a data element using quick sort requires many comparison operations. The combined processing time to perform the merge may be 4Tm(N/3). Thus the processing time for a quick sort is Tqs(N) where the processing time for the merge **800** is Tqs(N/9)+4Tm(N/3). Thus, the merge **800** may be performed in ⅙ the time as a quick sort operation. The above-described merge **800** may be applicable for merging 3, 9, 27, 81, or other power of 3 number of data sets. The above-described application may not be applicable for merging other numbers of data sets, such as 5, 7, 10, 12, or some other number.

FIG. **9** illustrates an example system **900** for partitioning data sets, arranged in accordance with at least some embodiments described herein. The system **900** may include shared memory **910**, a first processor **912**, a second processor **914**, a third processor **916**, a fourth processor **918**, a fifth processor **920**, and a sixth processor **922**. The shared memory **910** may include a data set A **930**, a data set B **932**, and a data set C **934**.

Each of the processors **912**, **914**, **916**, **918**, **920**, **922** may be configured to partition the data sets **930**, **932**, **934** in a single location. The processors **912**, **914**, **916**, **918**, **920**, **922** may perform a method similar to the method described above with respect to FIGS. **3**A-**3**C to partition the data sets **930**, **932**, **934**. In some embodiments, the processors **912**, **914**, **916**, **918**, **920**, **922** may be configured to partition the data sets **930**,

932, 934 in parallel. For example, each of the processors 912, 914, 916, 918, 920, 922 may be configured to determine a partition index for each of the data sets 930, 932, 934 in substantially overlapping time intervals. In these and other embodiments, the processors 912, 914, 916, 918, 920, 922 may partition the data sets 930, 932, 934 in six locations to create seven portions for each of the data sets 930, 932, 934.

In some embodiments, fewer than all of the processors 912, 914, 916, 918, 920, 922 may determine a partition index for each of the data sets 930, 932, 934. In these and other embodiments, two of the processors 912, 914, 916, 918, 920, 922 may collaborate to determine a partition index for each of the data sets 930, 932, 934. For example, the first processor 912 may perform a portion of the partitioning of the data sets 930, 932, 934 by determining when a reversal occurs. The first processor 912 may then determine the tail data elements of a first portion of the data sets 930, 932, 934 and the second processor 914 may determine the head data elements of a second portion of the data sets 930, 932, 934.

The shared memory 910 may include instructions for performing a partition of the data sets 930, 932, and 934. Each of the processors 912, 914, 916, 918, 920, 922 may access the instructions. In some embodiments, one of the processors 912, 914, 916, 918, 920, 922, such as the first processor 912 may be configured to coordinate the partitioning of the data sets among the processors 912, 914, 916, 918, 920, 922. In these and other embodiments, the first processor 912 may be configured to determine initial partition indexes for each of the processors 912, 914, 916, 918, 920, 922. Alternately or additionally, the first processor 912 may be configured to receive an indication of when the processors 912, 914, 916, 918, 920, 922 have completed the partitioning of the data sets 930, 932, 934. Alternately or additionally, the first processor 912 may coordinate processing of separate portions of the data sets 930, 932, 934 produced by the partitioning of the data sets 930, 932, 934 in a manner similar to or different from that described with respect to FIGS. 2A and 2B.

In some embodiments, the processors 912, 914, 916, 918, 920, 922 may contain cache memory that may be used in partitioning the data sets 930, 932, 934. The processors 912, 914, 916, 918, 920, 922 may contain various kinds of mutual-exclusion mechanisms for using computer resources exclusively. Alternately or additionally, each processor 912, 914, 916, 918, 920, 922 may contain mechanisms for solving inconsistencies that may occur when sharing memory among the processors 912, 914, 916, 918, 920, 922.

Below is example pseudo code that may be used to perform a merge operation of nine data sets as illustrated and described above with respect to FIG. 8 using a system similar to the system 900 described above.

The pseudo code may use various functions such as a QS( ), MERGE3( ), and DIV_MERGE3( ) function. The function QS( ) may be designed to sort a data set using a quick sort mechanism when passed the head position of the data set and the number of data elements in the data set. The function MERGE3( ) may perform a 3 way merge of data sets when passed the position information of the data sets, the position information for the merged data set, and the number of data elements. The function DIV_MERGE3( ) may divide 3 pair type aligned data sets based on the method described in FIGS. 3A-3C and merge the 3 pair type aligned data sets based on the division when passed the head and end positions of the 3 pair type aligned data sets and an initial partition index. Each of the nine data sets that are merged are assigned a parameter list, such as parameter list-1, parameter list-2, that represents the information about each of the nine data sets that is used by the functions to perform the merge function.

In some embodiments, the pseudo-code may use a coding library such as a multi-platform shared-memory parallel programming library for programming in C, C++, Fortran, and other programming languages. In particular, the code below may use the syntax of OpenMP. OpenMP may allow for ease in performing functions in parallel when the functions are not written specifically for parallel processing, such as the functions described above.

In particular, the parallel syntax, sections syntax, and section syntax of OpenMP may be used to describe parallel execution of functions. In OpenMP, the blocks enclosed in the parallel syntax are performed in parallel. Additionally, the blocks enclosed in the sections syntax are each assigned to a thread. The section syntax specifically expresses the function that each thread executes.

(a) Description of a Quick Sort Group

```
#pragma omp parallel {
#pragma omp sections  {
#pragma omp section    { /* start of a section */
--;
QS(parameter list-1);
...
}           end of a section
#pragma omp section {
--;
QS(parameter list-9);
...;
}
}                 /* end of sections */
}                   /* end of parallel */
```

(b) Description of the First Stage Merge Group

```
#pragma omp parallel {
#pragma omp sections  {
#pragma omp section    {
---;
MERGE3 (parameter list1);
---
}                   /* end of a section */
#pragma omp section {
---;
DIV_MERGE3(parameter list2);
--
}
#pragma omp section {
---;
DIV_MERGE3(parameter list3);
--
}
#pragma omp section {
--;
MERGE3(parameter list4);
---
}
#pragma omp section {
--;
DIV_MERGE3(parameter list5);
---
}
#pragma omp section {
---;
DIV_MERGE3 (parameter list6); --}
---
}
#pragma omp section {
MERGE3(parameter list7);
}
#pragma omp section {
--;
DIV_MERGE3(parameter list8);
--
}
```

-continued

```
#pragma omp section {
--;
DIV_MERGE3(parameter list9);
---
}          / end of section    */
}           /* end of sections */
}            /* end of parallel */
```

(c) The Description of the 2nd Stage Merge Group

```
#pragma omp parallel {
#pragma omp sections {
#pragma omp section {
--;
MERGE3 (parameter list-1);
--,
}
#pragma omp section {
--;
DIV_MERGE3(parameter list-2);
--
}
...
#pragma omp section {
--,
DIV_MERGE3(parameter list-9);
--
} /* end of sections */
}  /* end of parallel */
```

FIG. **10** illustrates another example system **1000** for partitioning data sets, arranged in accordance with at least some embodiments described herein. The system **1000** may include a connection network **1010**, a first processor **1020** coupled to a first private memory (PM) **1022**, a second processor **1030** coupled to a second PM **1032**, a third processor **1040** coupled to a third PM **1042**, a fourth processor **1050** coupled to a fourth PM **1052**, a communications processor **1060**, and a main PM **1070** that may include a data set A **1072**, a data set B **1074**, and a data set C **1076**. The system **1000** may also include a first connector **1024** coupled between the connection network **1010** and the first processor **1020** and the first PM **1022**, a second connector **1034** coupled between the connection network **1010** and the second processor **1030** and the second PM **1032**, a third connector **1044** coupled between the connection network **1010** and the third processor **1040** and the third PM **1042**, a fourth connector **1054** coupled between the connection network **1010** and the fourth processor **1050** and the fourth PM **1052**, a connector **1062** coupled between the connection network **1010** and the processor **1060** and the private memory **1070**.

The first connector **1024**, the second connector **1034**, the third connector **1046**, the fourth connector **1056** and the connector **1062** may be configured to enable the processors **1020**, **1030**, **1040**, **1050**, **1060** and the PMs **1022**, **1032**, **1042**, **1052**, **1070** respectively to communicate over the connection network **1010**.

Each of the processors **1020**, **1030**, **1040**, **1050** may be configured to communicate with the main PM **1070** over the connection network **1010** by way of the connectors **1024**, **1034**, **1046**, **1056** respectively. In some embodiments, the communications processor **1060** may facilitate communications between the main private memory **1070** and the processors **1020**, **1030**, **1040**, **1050**.

Each of the processors **1020**, **1030**, **1040**, **1050** may be configured to partition the data sets **1072**, **1074**, **1076** in a single location. The processors **1020**, **1030**, **1040**, **1050** may perform a method similar to the method described above with respect to FIGS. **3**A-**3**C to partition the data sets **1072**, **1074**, **1076**. In some embodiments, the processors **1020**, **1030**, **1040**, **1050** may be configured to partition the data sets **1072**, **1074**, **1076** in parallel. For example, each of the processors **1020**, **1030**, **1040**, **1050** may be configured to determine a partition index for each of the data sets **1072**, **1074**, **1076** in substantially overlapping time intervals. In these and other embodiments, the processors **1020**, **1030**, **1040**, **1050** may partition the data sets **1072**, **1074**, **1076** in four locations to create five portions. To partition the data sets **1072**, **1074**, **1076**, the processors **1020**, **1030**, **1040**, **1050** may each request and receive the data sets **1072**, **1074**, **1076** from the main PM **1070** over the connection network **1010**. Each of the processors **1020**, **1030**, **1040**, **1050** may store the data sets **1072**, **1074**, **1076** in their respectively coupled PM **1022**, **1032**, **1042**, **1052**. In some embodiments, the processors **1020**, **1030**, **1040**, **1050** may also receive instructions regarding partitioning the data sets **1072**, **1074**, **1076** from the main PM **1070** over the connection network **1010**. Alternately or additionally, each of the PMs **1022**, **1032**, **1042**, **1052** may contain instructions regarding partitioning the data sets **1072**, **1074**, **1076**.

In some embodiments, the communications processor **1060** may be configured to coordinate the partitioning of the data sets among the processors **1020**, **1030**, **1040**, **1050**. In these and other embodiments, the communications processor **1060** may be configured to determine initial partition indexes for each of the processors **1020**, **1030**, **1040**, **1050**. Alternately or additionally, the communications processor **1060** may be configured to receive an indication of when the processors **1020**, **1030**, **1040**, **1050** have completed the partitioning of the data sets **1072**, **1074**, **1076**. Alternately or additionally, the communications processor **1060** may coordinate processing of separate portions of the data sets **1072**, **1074**, **1076** produced by the partitioning of the data sets **1072**, **1074**, **1076** in a manner similar to or different from that described with respect to FIGS. **2**A and **2**B.

FIG. **11** illustrates another example system **1100** for partitioning data sets, arranged in accordance with at least some embodiments described herein. The system **1100** may include a switch network **1110**, a first processor **1120**, a second processor **1130**, a third processor **1140**, and a database **1150** that may include a data set A **1152**, a data set B **1154**, and a data set C **1156**.

Each of the processors **1120**, **1130**, **1140** may be configured to communicate with the database **1150** by way of the switch network **1110**. Each of the processors **1120**, **1130**, **1140** may be configured to partition the data sets **1152**, **1154**, **1156** in a single location. The processors **1120**, **1130**, **1140** may perform a method similar to the method described above with respect to FIGS. **3**A-**3**C to partition the data sets **1152**, **1154**, **1156**. The system **1100** may operate similar to the operation of the systems **900** and **1000** described herein with respect to FIGS. **9** and **10**.

The present disclosure is not to be limited in terms of the particular embodiments described herein, which are intended as illustrations of various aspects. Many modifications and variations can be made without departing from its spirit and scope, as will be apparent to those skilled in the art. Functionally equivalent methods and apparatuses within the scope of the disclosure, in addition to those enumerated herein, will be apparent to those skilled in the art from the foregoing descriptions. Such modifications and variations are intended to fall within the scope of the appended claims. The present disclosure is to be limited only by the terms of the appended claims, along with the full scope of equivalents to which such claims are entitled. It is to be understood that the present

disclosure is not limited to particular methods, reagents, compounds compositions or biological systems, which can, of course, vary. It is also to be understood that the terminology used herein is for the purpose of describing particular embodiments only, and is not intended to be limiting.

With respect to the use of substantially any plural and/or singular terms herein, those having skill in the art can translate from the plural to the singular and/or from the singular to the plural as is appropriate to the context and/or application. The various singular/plural permutations may be expressly set forth herein for sake of clarity.

It will be understood by those within the art that, in general, terms used herein, and especially in the appended claims (e.g., bodies of the appended claims) are generally intended as "open" terms (e.g., the term "including" should be interpreted as "including but not limited to," the term "having" should be interpreted as "having at least," the term "includes" should be interpreted as "includes but is not limited to," etc.). It will be further understood by those within the art that if a specific number of an introduced claim recitation is intended, such an intent will be explicitly recited in the claim, and in the absence of such recitation no such intent is present. For example, as an aid to understanding, the following appended claims may contain usage of the introductory phrases "at least one" and "one or more" to introduce claim recitations. However, the use of such phrases should not be construed to imply that the introduction of a claim recitation by the indefinite articles "a" or "an" limits any particular claim containing such introduced claim recitation to embodiments containing only one such recitation, even when the same claim includes the introductory phrases "one or more" or "at least one" and indefinite articles such as "a" or "an" (e.g., "a" and/or "an" should be interpreted to mean "at least one" or "one or more"); the same holds true for the use of definite articles used to introduce claim recitations. In addition, even if a specific number of an introduced claim recitation is explicitly recited, those skilled in the art will recognize that such recitation should be interpreted to mean at least the recited number (e.g., the bare recitation of "two recitations," without other modifiers, means at least two recitations, or two or more recitations). Furthermore, in those instances where a convention analogous to "at least one of A, B, and C, etc." is used, in general such a construction is intended in the sense one having skill in the art would understand the convention (e.g., "a system having at least one of A, B, and C" would include but not be limited to systems that have A alone, B alone, C alone, A and B together, A and C together, B and C together, and/or A, B, and C together, etc.). In those instances where a convention analogous to "at least one of A, B, or C, etc." is used, in general such a construction is intended in the sense one having skill in the art would understand the convention (e.g., "a system having at least one of A, B, or C" would include but not be limited to systems that have A alone, B alone, C alone, A and B together, A and C together, B and C together, and/or A, B, and C together, etc.). It will be further understood by those within the art that virtually any disjunctive word and/or phrase presenting two or more alternative terms, whether in the description, claims, or drawings, should be understood to contemplate the possibilities of including one of the terms, either of the terms, or both terms. For example, the phrase "A or B" will be understood to include the possibilities of "A" or "B" or "A and B."

In addition, where features or aspects of the disclosure are described in terms of Markush groups, those skilled in the art will recognize that the disclosure is also thereby described in terms of any individual member or subgroup of members of the Markush group.

As will be understood by one skilled in the art, for any and all purposes, such as in terms of providing a written description, all ranges disclosed herein also encompass any and all possible sub ranges and combinations of sub ranges thereof. Any listed range can be easily recognized as sufficiently describing and enabling the same range being broken down into at least equal halves, thirds, quarters, fifths, tenths, etc. As a non-limiting example, each range discussed herein can be readily broken down into a lower third, middle third and upper third, etc. As will also be understood by one skilled in the art all language such as "up to," "at least," and the like include the number recited and refer to ranges which can be subsequently broken down into sub ranges as discussed above. Finally, as will be understood by one skilled in the art, a range includes each individual member. Thus, for example, a group having 1-3 cells refers to groups having 1, 2, or 3 cells. Similarly, a group having 1-5 cells refers to groups having 1, 2, 3, 4, or 5 cells, and so forth.

From the foregoing, it will be appreciated that various embodiments of the present disclosure have been described herein for purposes of illustration, and that various modifications may be made without departing from the scope and spirit of the present disclosure. Accordingly, the various embodiments disclosed herein are not intended to be limiting, with the true scope and spirit being indicated by the following claims.

The invention claimed is:

1. A method for a computing device, the method comprising:

locating, using the computing device, a first partition index for each of three sorted data sets, each sorted data set including a plurality of indexed data values, wherein each first partition index identifies an index location to partition its corresponding sorted data set into a first portion of data values and a second portion of data values, wherein each data value in each of the first portions of each of the three sorted data sets has a greater magnitude than each data value in each of the second portions of each of the three sorted data sets, wherein locating the first partition index for each of the sorted data sets comprises:

selecting an initial partition index for each sorted data set;

comparing data values at the initial partition index for each sorted data set to identify a first sorted data set of the three sorted data sets with a highest data value of the compared data values and a second sorted data set of the three sorted data sets with a lowest data value of the compared data values;

adjusting the initial partition index for the first sorted data set identified with the highest data value to an adjusted initial partition index with a lower data value;

adjusting the initial partition index for the second sorted data set identified with the lowest data value to an adjusted initial partition index with a higher data value; and

comparing data values at the adjusted initial partition indexes of the first and second sorted data sets with the data value at the initial partition index of a third sorted data set of the three sorted data sets without an adjusted initial partition index.

2. The method of claim **1**, wherein if the initial partition index of the second sorted data set with the lowest data value is adjusted by incrementing the initial partition index, the initial partition index of the first sorted data set with the highest data value is adjusted by decrementing the initial partition index.

3. A method for a computing device, the method comprising:

locating, using the computing device, a first partition index for each of three sorted data sets, each sorted data set including a plurality of indexed data values, wherein each first partition index identifies an index location to partition its corresponding sorted data set into a first portion of data values and a second portion of data values, wherein each data value in each of the first portions of each of the three sorted data sets has a greater magnitude than each data value in each of the second portions of each of the three sorted data sets, wherein locating the first partition index for each of the sorted data sets comprises:

selecting an initial partition index for each sorted data set;

comparing data values at the initial partition index for each sorted data set to identify a highest data value and a lowest data value;

adjusting the initial partition index for the sorted data set with the highest data value and for the sorted data set with the lowest data value;

comparing data values at the adjusted initial partition indexes of the sorted data sets having the adjusted initial partition indexes and the data value at the initial partition index of the sorted data set without an adjusted initial partition index; and

iterating over:

comparing data values at a current partition index for each sorted data set to identify a highest data value and a lowest data value at the current partition index; and

adjusting the current partition index for the sorted data sets with the highest and lowest data values at the current partition index,

until, during a single iteration, one of the sorted data sets that was previously determined to have the lowest data value changes to have the highest value and one of the sorted data sets that was previously determined to have the highest data value changes to have the lowest data value.

4. The method of claim 3, wherein locating the first partition index comprises:

utilizing the current partition indexes of the sorted data sets after the single iteration to determine the first partition index for each of the sorted data sets.

5. The method of claim 1, wherein each of the sorted data sets includes at least one data value that overlaps ranges of data values in one of the other sorted data sets.

6. The method of claim 1, wherein the initial partition index includes a first initial partition index and the index location includes a first index location, the method further comprising:

locating, using the computing device, a second partition index for each sorted data set, each second partition index identifying a second index location to partition the corresponding sorted data set so that each sorted data set forms a third portion of data values, wherein each data value in each of the third portions of the sorted data sets has a greater magnitude than each data value in each of the first and second portions of the sorted data sets, wherein locating the second partition index for each of the sorted data sets comprises:

selecting a second initial partition index for each sorted data set;

comparing data values at the second initial partition index for each sorted data set to identify a highest data value and a lowest data value;

adjusting the second initial partition index for the sorted data set with the highest data value and for the sorted data set with the lowest data value; and

comparing data values at the adjusted second initial partition indexes of the sorted data sets having the adjusted second initial partition indexes and the data value at the second initial partition index of the sorted data set without an adjusted second initial partition index.

7. The method of claim 6, wherein locating each first partition index is independent of locating each second partition index.

8. The method of claim 6, wherein processing time intervals to locate the first and second partition indexes occur in substantially overlapping time intervals.

9. The method of claim 6, wherein locating the first partition indexes is performed by a first processor and locating the second partition indexes is performed by a second processor.

10. A method comprising:

partitioning, by a computing device, three sorted data sets that each include a plurality of indexed data values into first and second portions so that each data value in each of the first portions of each of the three sorted data sets has a greater magnitude than each data value in each of the second portions of each of the three sorted data sets, wherein partitioning the three sorted data sets comprises:

selecting an initial partition index for each of the three sorted data sets;

comparing data values at the initial partition index for each of the three sorted data sets to identify a first sorted data set of the three sorted data sets with a highest data value of the compared data values and a second sorted data set of the three sorted data sets with a lowest data value of the compared data values;

adjusting the initial partition index for the first sorted data set identified with the highest data value to an adjusted initial partition index with a lower data value;

adjusting the initial partition index for the second sorted data set identified with the lowest data value to an adjusted initial partition index with a higher data value; and

comparing data values at the adjusted initial partition indexes of the first and second sorted data sets with the data value at the initial partition index of a third sorted data set of the three sorted data sets without an adjusted partition index;

processing the first portions of the sorted data sets; and

processing the second portions of the sorted data sets.

11. The method of claim 10, wherein processing the first portions of the sorted data sets is independent of processing the second portions of the sorted data sets.

12. The method of claim 10, wherein a first processing time interval to process the first portions of the sorted data sets and a second processing time interval to process the second portions of the sorted data sets occur in substantially overlapping time intervals.

13. The method of claim 10, wherein processing the first portions of the sorted data sets is performed by a first processor and processing the second portions of the sorted data sets is performed by a second processor.

14. The method of claim 10, wherein the first portions of the sorted data sets are processed by a first processor in a cloud computing environment and the second portions of the sorted data sets are processed by a second processor in the cloud computing environment.

15. The method of claim 14, wherein the three sorted data sets are partitioned in the cloud computing environment.

16. A system to process data, the system comprising:
a first processor configured to process a first portion of each of three sorted data sets; and
a second processor configured to process a second portion of each of the three sorted data sets,
wherein the three sorted data sets each include a plurality of indexed data values and the three sorted data sets are partitioned into the first and second portions so that each data value in each of the first portions of each of the three sorted data sets has a greater magnitude than each data value in each of the second portions of each of the three sorted data sets, wherein the first and second portions are partitioned by:
selection of an initial partition index for each of the three sorted data sets;
comparison of data values at the initial partition index for each of the three sorted data sets to identify a first sorted data set of the three sorted data sets with a highest data value of the compared data values and a second sorted data set of the three sorted data sets with a lowest data value of the compared data values;
adjustment of the initial partition index for the first sorted data set identified with the highest data value to an adjusted initial partition index with a lower data value;
adjustment of the initial partition index for the second sorted data set identified with the lowest data value to an adjusted initial partition index with a higher data value; and
comparison of data values at the adjusted initial partition indexes of the first and second sorted data sets with the data value at the initial partition index of a third sorted data set of the three sorted data sets without an adjusted partition index.

17. The system of claim 16, wherein the first processor is configured to process the first portions of the three sorted data sets substantially independent from the second processor and during substantially overlapping time intervals in which the second processor processes the second portions of the three sorted data sets.

18. The system of claim 16, wherein the first and second processors are each a separate core of a multi-core processor.

19. The system of claim 16, wherein:
the first processor or the second processor is configured to partition each of the three sorted data sets into the first and second portions; or
a third processor is configured to partition each of the three sorted data sets into the first and second portions.

20. The system of claim 16, wherein the first and second processors are configured to either share a memory unit or access separate memory units.

21. The system of claim 16, further comprising a third processor configured to process a third portion of each of the three sorted data sets, wherein each data value in each of the third portions of the sorted data sets has a greater magnitude than each data value in each of the first and second portions of the sorted data sets.

22. A method for a computing device, the method comprising:
locating, using the computing device, a first partition index for each of three sorted data sets, each sorted data set including a plurality of indexed data values, wherein each first partition index identifies an index location to partition the corresponding sorted data set into a first portion of data values and a second portion of data values, wherein each data value in each of the first portions of the sorted data sets has a greater magnitude than each data value in each of the second portions of the sorted data sets, wherein locating the first partition index for each of the sorted data sets comprises:
selecting an initial partition index for each sorted data set;
comparing data values at the initial partition index for each sorted data set to identify a highest data value and a lowest data value;
adjusting the initial partition index of the sorted data set identified with the lowest data value to a partition index with a higher value;
adjusting the initial partition index of the sorted data set identified with the highest data value to a partition index with a lower value;
comparing data values at the adjusted initial partition indexes of the sorted data sets having the adjusted initial partition indexes and the data value at the initial partition index of the sorted data set without an adjusted initial partition index;
setting the adjusted initial partition indexes of the sorted data sets having the adjusted initial partition indexes as current partition indexes for these sorted data sets;
setting the initial partition index of the sorted data set without an adjusted initial partition index as a current partition index for this sorted data set;
iterating over:
comparing data values at the current partition index for each sorted data set to identify a highest data value and a lowest data value at the current partition indexes; and
adjusting the current partition index for the sorted data sets with the highest and lowest data values at the current partition indexes,
until, during a single iteration, a sorted data set that was previously determined to have the lowest data value changes to have the highest data value and a sorted data set that was previously determined to have the highest data value changes to have the lowest data value;
after the sorted data set that was previously determined to have the lowest data value changes to have the highest data value and the sorted data set that was previously determined to have the highest data value changes to have the lowest data value, determining first and second candidate first partition indexes based on data values at the current partition indexes and data values at previous partition indexes;
comparing the first and second candidate first partition indexes to locate the first partition index for one of the sorted data sets; and
locating the first partition indexes for the sorted data sets that are not the one of the sorted data sets based on the first partition index for the one of the sorted data sets;
partitioning the sorted data sets into the first and second portions based on the located first partition index for each of the sorted data sets;
processing the first portions of the sorted data sets; and
processing the second portions of the sorted data sets.

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